Map Construction and Comparison Using Local Structure

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> 6 February 2014 SAMSI Workshop





Road networks:

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Networks:

- Road networks: GPS trajectories of cars.
- Hiking paths: GPS paths of people.
- Migration paths: GPS on animals.
- Fillaments of galaxies: point cloud data.
- Biological systems: Marron's brain vessels.
- Hurricane paths: historical paths.
- Networks: path-constrained trajectories.



Road Network Representation

Definition (Road Network)

A road network G = (V, E) is an embedded 1-complex. Edges E represent the *roads* or *streets* and the vertices V represent the *intersections*.



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The Data GPS Trajectories



Part I

The *Link Lengh* of a path is the number of edges in the path.

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Map Construction and Comparison

The Data GPS Trajectories



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Map Construction and Comparison

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Trajectory Data

Athens Trajectories from School Bus Data



The Reconstruction Problem

Problem Statement

Given a set of (constrained) trajectories, extract the underlying geometric graph structure, in particular:

- Intersections (Vertices of degree > 2)
- Streets (Edges: PL paths between vertices)

Existing Reconstructing Algorithms

http://www.mapconstruction.org

[ACCGGM-12] M. Aanjaneya, F. Chazal, D. Chen, M. Glisse, L. Guibas, and D. Morozov. Metric Graph Reconstruction from Noisy Data, 2012.

[BE-12a] J. Biagioni, J. Eriksson. Map Inference in the Face of Noise and Disparity. SIGSPATIAL, 2012.

[KP-12] S. Karagiorgou and D. Pfoser. On Vehicle Tracking Data-Based Road Network Generation. SIGSPATIAL, 2012.

[AW-12] M. Ahmed, C. Wenk. Constructing Street Networks from GPS Trajectories. European Symp. on Algorithms, 2012.

[OTHER] There were too many algorithms to cite here ...

Incremental Fréchet Matching [AW]



Assumption A1

Trajectories are Close to Paths



Incremental Fréchet Matching [AW]



Incremental Fréchet Matching [AW]

Assumptions

- A1: $d_F(p, \hat{p}) \leq \frac{\varepsilon}{2}$ for all trajectories \hat{p} .
- **2** A2: If two paths in G come close enough, they intersect.

Part I

Assumption A2

Close Paths Intersect



Assumption A2

Close Paths Intersect



Incremental Fréchet Matching [AW]

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Part I

Incremental Fréchet Matching [AW]

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- A1: $d_F(p, \hat{p}) \leq \frac{\varepsilon}{2}$ for all trajectories \hat{p} .
- **2** A2: If two paths in G come close enough, they intersect.
- Solution A3: For each street s and a ball B far enough away from the endpoints of s, s ∩ B is two points.

Assumption A3

Roads Are Nice



Assumption A3

For each street s and a ball B far enough away from the endpoints of s, $s \cap B$ is two points.

Incremental Fréchet Matching [AW]

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Incremental Fréchet Algorithm [AW]

For each trajectory *t*:

- Find the portion of the map closest to t.
- Add edges if necessary.
- **3** Simplify:

Incremental Fréchet Algorithm [AW]



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Computing the Fréchet Distance



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Finding Closest Match ...

Partial Fréchet Distance



Let grey region have weight 1 and white region weight 0. Find the shortest path from one end to the other under this metric.

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Adding Edges



Part I

Adding Edges



Part I

Incremental Fréchet Algorithm [AW]

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Incremental Fréchet Algorithm [AW]

For each trajectory *t*:

- Find the portion of the map closest to t.
- Add edges if necessary.
- Simplify: Min-link curve simplification algorithm.

Dataset 1: Berlin

- 27,188 GPS trajectories from school buses.
- D = 5.2km × 6.1km.
- Trajectories have 7 observations on average.

Dataset 1: Berlin



Dataset 1: Berlin



Dataset 1: Berlin



Dataset 1: Berlin



Dataset 2: Chicago

- 889 GPS trajectories from university buses.
- D = 7 km x 4.5 km
- Trajectories: 100-300 samples (avg. 3.22 km).

Dataset 2: Chicago



Dataset 2: Chicago



Dataset 2: Chicago



Dataset 2: Chicago



Dataset 2: Chicago





Different Approaches

[BE-12b] J. Biagioni, J. Eriksson. James Biagioni and Jakob Eriksson. Inferring Road Maps from GPS Traces: Survey and Comparative Evaluation. In 91st Annual Meeting of the Transportation Research Board, 2012.

[AFHW-14] M. Ahmed, B. Fasy, K. Hickmann, C. Wenk. Path-Based Distance for Street Map Comparison. Arxiv.

[AFW-14] M. Ahmed, B. Fasy, C. Wenk. Local Homology Based Distance Between Maps. In Submission.










Local Topology Approach [AFW]



We will use the embedding and the local homology to create a *local distance signature*.

Local Distance Signature

Finding the Local Persistence Diagram



 $D \subset \mathbb{R}^2$ is the compact domain. $G_1 \subset D$ is the road network. r > 0 is the scale. $x \in D$.

$LG_1(x,0) = (G_1 \cap B_r(x))/\partial B_r(x)$

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Local Distance Signature

Finding the Local Persistence Diagram



Part II

Local Distance Signature

Computing the Local Distance



Local Distance Signature

Computing the Local Distance



Local Topology Based Distance

Definition (Local Homology Distance)

$$d^{LH}(G_1,G_2)=\frac{1}{|D|} \qquad \int_D lhd(x,r)\,dx$$



$$d^{LH}(G_1, G_2) = \frac{1}{|D|} \int_{r_0}^{r_1} \int_D Ihd(x, r) \, dx \, dr$$

Definition (Local Homology Distance)

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Variants

- D is a rectangular domain.
- **2** $D = (G_1 \cup G_2) + B_{\delta}$
- 3 $r_0 = r_1$ (fixed radius)
- $I r_0 = 0 \implies metric$

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Dataset 3: Athens

www.openstreetmaps.org

- 129 GPS trajectories from school buses.
- D = 2.6 km x 6 km
- Two reconstruction algorithms: [BE-12a] and [KP-12].
- Trajectories: 13-47 samples

Dataset 3: Athens

www.openstreetmaps.org



Dataset 3: Athens

www.openstreetmaps.org



Dataset 3: Athens

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Thanks Fabrizio and Jisu!

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Results

Athens: Comparing Two Different Reconstructions



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Athens: Comparing Two Different Reconstructions





The Bootstrap

... When the Ground Truth is Unknown

- *n* input trajectories S_n .
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$$T_j = d^{LH}(\widehat{G}, \widehat{G}_j^*).$$

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Ongoing

- Apply this to different construction algorithms versus the ground truth to rank algorithms.
- Implementation of distance measure.
- Improve theoretical guarantees.
- Input Model: other noise models?
- Output Model: road category, direction, intersection regions, ...

Summary





Map Construction

[AW-12] Constructing Street Networks from GPS Trajectories.

Map Comparison

[AFW-14] Local Homology Based Distance Between Maps.

More References

[AG-95] H. Alt, M. Godau. Computing the Frchet Distance Between Two Polygonal Curves. IJCGA, 1995.

[BBW-09] K. Buchin, M. Buchin, Y. Wang. Exact Algorithms for Partial Curve Matching via the Fréchet Distance. SODA 2009.

[CW-10] A. F. Cook IV, C. Wenk. Geodesic Frchet Distance Inside a Simple Polygon. ACM TALG 7(1), 2010.

Thank You!

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