Statistical Computing (36-350)
Lecture 18: Optimization II: Stochastic and Constrained Optimization

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Stochastic Optimization
Constraints and Penalties

Agenda

- Stochastic optimization methods
- Constraints and penalties

OPTIONAL READING (big picture): Francis Spufford, *Red Plenty*;
Herbert Simon, *The Sciences of the Artificial*

OPTIONAL READING (close up): Bottou and Bosquet, “The Tradeoffs of Large Scale Learning”
Problems with Big Data

Typical statistical objective function, mean-squared error:

\[ f(\theta) = \frac{1}{n} \sum_{i=1}^{n} (y_i - m(x_i, \theta))^2 \]

Getting a value of \( f \) is \( O(n) \), \( \nabla f \) is \( O(np) \), \( H \) is \( O(np^2) \)

worse still if \( m \) slows down with \( n \)

Not bad when \( n = 100 \) or even \( n = 10^4 \), but if \( n = 10^9 \) or \( n = 10^{12} \) we don’t even know which way to move
"Oh sure, going in that direction will totally minimize the objective function" — Sarcastic Gradient Descent.

20 Jul 12
Sampling, the Alternative to Sarcastic Gradient Descent

Pick one data point $I$ at random (uniform on 1 : $n$) Loss there, $(y_I - m(x_I, \theta))^2$, is random, but

\[ \mathbb{E} [(y_I - m(x_I, \theta))^2] = f(\theta) \]

Generally, if $f(\theta) = n^{-1} \sum_{i=1}^{n} f_i(\theta)$ and $f_i$ are well-behaved,

\[ \mathbb{E} [f_I(\theta)] = f(\theta) \]
\[ \mathbb{E} [\nabla f_I(\theta)] = \nabla f(\theta) \]
\[ \mathbb{E} [\nabla^2 f_I(\theta)] = \nabla^2 f(\theta) \]

∴ Don’t optimize with all the data, optimize with random samples
Stochastic Gradient Descent

Draw lots of one-point samples, let their noise cancel out:

1. Start with initial guess $\theta$, learning rate $\eta$
2. While ((not too tired) and (making adequate progress))
   1. At $t^{th}$ iteration, pick random $I$ uniformly
   2. Set $\theta \leftarrow \theta - t^{-1} \eta \nabla f_I(\theta)$
3. Return final $\theta$

Shrinking step-size by $1/t$ ensures noise in each gradient dies down
(Variants: put points in some random order, only check progress after going over each point once, adjust $1/t$ rate, average a couple of random data points, etc.)

The sample code from the midterm works, though it could be made more efficient
Stochastic Newton’s Method

a.k.a. 2nd order stochastic gradient descent

1. Start with initial guess $\theta$
2. While ((not too tired) and (making adequate progress))
   1. At $t^{th}$ iteration, pick uniformly-random $I$
   2. $\theta \leftarrow \theta - t^{-1}H_{I}^{-1}(\theta)\n\n+ all the Newton-ish tricks to avoid having to recompute the Hessian
Stochastic Gradient Methods

Pros:
- Each iteration is fast (and constant in $n$)
- Never need to hold all data in memory
- Does converge eventually

Cons:
- Noise *does* reduce precision — more iterations to get within $\epsilon$ of optimum than non-stochastic GD or Newton

Often low computational cost to get within *statistical* error of the optimum
Simulated Annealing

Use Metropolis to sample from a density $\propto e^{-f(\theta)/T}$

Samples will tend to be near small values of $f$

Keep lowering $T$ as we go along ("cooling", "annealing")

1. Set initial $\theta$, $T > 0$
2. While ((not too tired) and (making adequate progress))
   1. Proposal: $Z \leftarrow r(\cdot|\theta)$ (e.g., Gaussian noise)
   2. Draw $U \sim \text{Unif}(0,1)$
   3. Acceptance: If $U < e^{-\frac{f(Z) - f(\theta)}{T}}$ then $\theta \leftarrow Z$
   4. Reduce $T$ a little
3. Return final $\theta$

Always moves to lower values of $f$, sometimes moves to higher
No derivatives, works for discrete problems, few guarantees
Maximizing a multinomial likelihood

I roll dice \( n \) times; \( n_1, \ldots, n_6 \) count the outcomes

Likelihood and log-likelihood:

\[
L(\theta_1, \theta_2, \theta_3, \theta_4, \theta_5, \theta_6) = \frac{n!}{n_1!n_2!n_3!n_4!n_5!n_6!} \prod_{i=1}^{6} \theta_i^{n_i}
\]

\[
\ell(\theta_1, \theta_2, \theta_3, \theta_4, \theta_5, \theta_6) = \log \left( \frac{n!}{n_1!n_2!n_3!n_4!n_5!n_6!} \right) + \sum_{i=1}^{6} n_i \log \theta_i
\]

Optimize by taking the derivative and setting to zero:

\[
\frac{\partial \ell}{\partial \theta_1} = \frac{n_1}{\theta_1} = 0
\]

\[
\therefore \theta_1 = \infty
\]

or \( n_1 = 0 \)
We forgot that $\sum_{i=1}^{6} \theta_i = 1$

We could use the constraint to eliminate one of the variables

$$\theta_6 = 1 - \sum_{i=1}^{5} \theta_i$$

Then solve the equations

$$\frac{\partial \mathcal{L}}{\partial \theta_i} = \frac{n_1}{\theta_i} - \frac{n_6}{1 - \sum_{j=1}^{5} \theta_j} = 0$$

BUT eliminating a variable with the constraint is usually messy
Lagrange Multipliers

\[ g(\theta) = c \iff g(\theta) - c = 0 \]

Lagrangian:

\[ \mathcal{L}(\theta, \lambda) = f(\theta) - \lambda(g(\theta) - c) \]

=\( f \) when the constraint is satisfied
Now do \textit{unconstrained} minimization over \( \theta \) and \( \lambda \):

\[ \nabla_{\theta} \mathcal{L} \big|_{\theta^*, \lambda^*} = \nabla f(\theta^*) - \lambda^* \nabla g(\theta^*) = 0 \]

\[ \frac{\partial \mathcal{L}}{\partial \lambda} \bigg|_{\theta^*, \lambda^*} = g(\theta^*) - c = 0 \]

optimizing \textbf{Lagrange multiplier} \( \lambda \) enforces constraint
More constraints, more multipliers
Try the dice again:

\[
\mathcal{L} = \log \frac{n!}{\prod_i n_i!} + \sum_{i=1}^{6} n_i \log(\theta_i) - \lambda \left( \sum_{i=1}^{6} \theta_i - 1 \right)
\]

\[
\frac{\partial \mathcal{L}}{\partial \theta_i} \bigg|_{\theta_i=\theta_i^*} = \frac{n_i}{\theta_i^*} - \lambda^* = 0
\]

\[
\frac{n_i}{\lambda^*} = \theta_i^*
\]

\[
\sum_{i=1}^{6} \frac{n_i}{\lambda^*} = \sum_{i=1}^{6} \theta_i^* = 1
\]

\[
\lambda^* = \sum_{i=1}^{6} n_i \Rightarrow \theta_i^* = \frac{n_i}{\sum_{i=1}^{6} n_i}
\]
Thinking About the Lagrange Multipliers

Constrained minimized is generally higher than the unconstrained
Changing the constraint level $c$ changes $\theta^*, f(\theta^*)$

$$\frac{\partial f(\theta^*)}{\partial c} = \frac{\partial L(\theta^*, \lambda^*)}{\partial c}$$

$$= [\nabla f(\theta^*) - \lambda^* \nabla g(\theta^*)] \frac{\partial \theta^*}{\partial c} - [g(\theta^*) - c] \frac{\partial \lambda^*}{\partial c} + \lambda^* = \lambda^*$$

$\lambda^*$ = Rate of change in optimal value as the constraint is relaxed
$\lambda^*$ = “Shadow price”: How much would you pay for minute change in the level of the constraint
What about an *inequality* constraint?

\[ h(\theta) \leq d \iff h(\theta) - d \leq 0 \]

The region where the constraint is satisfied is the **feasible set**

*Roughly* two cases:

1. Unconstrained optimum is inside the feasible set \( \Rightarrow \) constraint is **inactive**
2. Optimum is outside feasible set; constraint is **active**, **binds** or **bites**; *constrained* optimum is usually on the boundary

Add a Lagrange multiplier; \( \lambda \neq 0 \iff \) constraint binds
Older than computer programming...

Optimize $f(\theta)$ subject to $g(\theta) = c$ and $h(\theta) \leq d$

“Give us the best deal on $f$, keeping in mind that we’ve only got $d$ to spend, and the books have to balance”

Linear programming (Kantorovich, 1938)

- $f, h$ both linear in $\theta$
- $\theta^*$ always at a corner of the feasible set
Back to the Factory

Constraints:

\[ 40(\text{cars}) + 60(\text{trucks}) \leq 1600 \]
\[ 1(\text{cars}) + 3(\text{trucks}) \leq 70 \]

Revenue: $13k/car, $27k/truck

The feasible region:
Back to the Factory

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The feasible region:
Barrier Methods

(a.k.a. “interior point”, “central path”, etc.)
Having constraints switch on and off abruptly is annoying especially with gradient methods
Fix $\mu > 0$ and try minimizing

$$f(\theta) - \mu \log(d - h(\theta))$$

“pushes away” from the barrier — more and more weakly as $\mu \to 0$

1. Initial $\theta$ in feasible set, initial $\mu$
2. While ((not too tired) and (making adequate progress))
   1. Minimize $f(\theta) - \mu \log(d - h(\theta))$
   2. Reduce $\mu$
3. Return final $\theta$
argmin_{\theta} f(\theta) \iff argmin_{\theta, \lambda} f(\theta) - \lambda(h(\theta) - d)

\text{Constraint level } d \text{ doesn’t matter for doing the second minimization over } \theta

\text{Constrained optimization} \iff \text{Penalized optimization}

\text{Constraint level } d \iff \text{Penalty factor } \lambda
Minimize MSE of linear function $\beta \cdot x$: ordinary least squares regression
penalty on length of coefficient vector, $\sum \beta_j^2$: ridge regression
stabilizes estimate when data are noisy, $p > n$, collinear
penalty on sum of coefficients, $\sum |\beta_j|$: lasso
stability + drive small coefficients to 0 ("sparsity")
Minimize MSE of function + penalty on curvature: spline
fit smooth regressions w/o assuming specific form
Smoothing over time, space, other relations
e.g., social or genetic ties
 Usually decide on penalty factor/constraint level by trying to predict out of sample
In penalty form, just choose $\lambda$ and modify your objective function.

`constrOptim` implements the barrier method.

Try this:

```r
factory <- matrix(c(40,1,60,3), nrow=2,
                  dimnames=list(c("labor","steel"),c("car","truck")))
available <- c(1600,70); names(available) <- rownames(factory)
prices <- c(car=13, truck=27)
revenue <- function(output) { return(-output %*% prices) }
plan <- constrOptim(theta=c(5,5), f=revenue, grad=NULL,
                     ui=-factory, ci=-available, method="Nelder-Mead")
plan$par
```

only works with constraints like $u \theta \geq c$, so minus signs.
Summary

- Stochastic optimization methods use probability in the search
  - Stochastic gradient descent samples the data; gives up precision for speed
  - Simulated annealing randomly moves against the objective function; escapes local minima

- Constraints are usually part of optimization
  - Constrained optimum generally at the boundary of feasible set
  - Lagrange multipliers turn constrained problems into unconstrained ones
  - Multipliers are prices: trade-off between tightening constraint and worsening optimal value