

# **36-720: Graphical Models**

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- Collapsibility and Simpson's Paradox
- Decomposability
- Examples

One nice feature of graphical models is that they lead immediately to two important ideas in simplifying log-linear models. They are

- Decomposable models
- Collapsible models

Decomposable models are graphical models for which closed-form MLE's exist.

Collapsibility is closely related to (the non-occurrence of) Simpson's paradox like phenomena.

## Collapsibility and Simpson's Paradox

Consider a hierarchical, graphical log-linear model  $\mathcal{M}$  and partition the factors in  $\mathcal{M}$  into three disjoint groups,  $A$ ,  $B$ , and  $C$ . *Examples:*

- $\mathcal{M} = [12][23]$ ,  $A = \{1\}$ ,  $B = \{3\}$ ,  $C = \{2\}$
- $\mathcal{M} = [123][145]$ ,  $A = \{2, 3\}$ ,  $B = \{4, 5\}$ ,  $C = \{1\}$
- $\mathcal{M} = [126][234][46][56]$ ,  $A = \{1, 2, 6\}$ ,  $B = \{3, 5\}$ ,  $C = \{4\}$

Often we are interested in the relationship between  $A$  and  $C$

- Conditioning on  $B$ , vs.
- Summing (collapsing) over  $B$ .

Simpson's paradox happens when

$$P(A|C, B) \neq P(A|C)$$

[e.g. if  $A$  and  $C$  are single factors,  $OR(A, C) \neq OR(A, C|B)$ .]

We will say “the conditional relationship  $A|C$  is collapsible over  $B$ ” if

$$P(A|C, B) = P(A|C)$$

Clearly this can happen if and only if

$$B \perp\!\!\!\perp A \mid C$$

Since  $\mathcal{M}$  is a graphical model, the *global Markov property* tells us this can happen iff

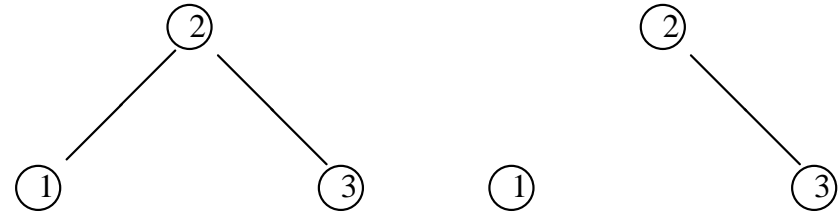
*$C$  separates  $\mathcal{M}$  into disjoint subgraphs containing  $A$  and  $B$ .*

and when we look at the graph we see this happens iff

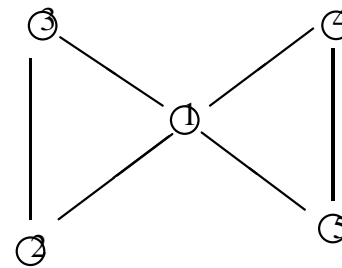
*Every path from a variable in  $A$  to a variable in  $B$  goes through at least one variable in  $C$ .*

Thus we have a simple criterion for the collapsibility of relationships between groups of variables, over a “stratification” variable. *Examples:*

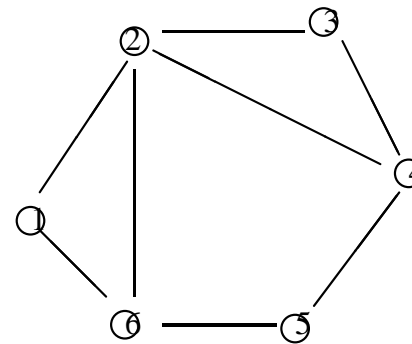
$\mathcal{M} = [12][23]$ ,  $A = \{1\}$ ,  $B = \{3\}$ ,  
 $C = \{2\}$



$\mathcal{M} = [123][145]$ ,  $A = \{2, 3\}$ ,  $B =$   
 $\{4, 5\}$ ,  $C = \{1\}$



$\mathcal{M} = [126][234][45][56]$ ,  $A =$   
 $\{1, 2, 6\}$ ,  $B = \{3, 5\}$ ,  $C = \{4\}$



## An Equivalent Form of Collapsibility

Sometimes we are not so much interested in the relationship between variables in  $A$  and those in  $C$ , but rather in *all* the relationships in  $A$ , after collapsing over  $B$ .

Let  $A$  be a subset of variables in  $\mathcal{M}$  and let  $B = \mathcal{M} \setminus A$ .

We say  $\mathcal{M}$  is *collapsible onto*  $A$  iff the models

- $P(A)$
- $P(A|B)$

have the same graph [namely, the graph of  $\mathcal{M}$  with the variables in  $B$  deleted].

I.e.,  $P(A)$  has exactly the same conditional independence relationships as  $P(A|B)$ .

*The two forms of collapsibility are equivalent:*

- Let  $A = A' \cup C'$  [disjoint union]. Clearly, if  $\mathcal{M}$  is collapsible onto  $A$ , then  $P(A'|C') = P(A'|C', B)$  since these are determined by  $P(A)$  and  $P(A|B)$ .
- Conversely if  $A'|C'$  is collapsible over  $B$ , then the table can be collapsed onto  $A = A' \cup B'$ . This follows from Theorem 2.3 and Corollary 2.5 of Asmussen & Edwards (1983, *Bmka*), which also establishes the criterion on the next slide.

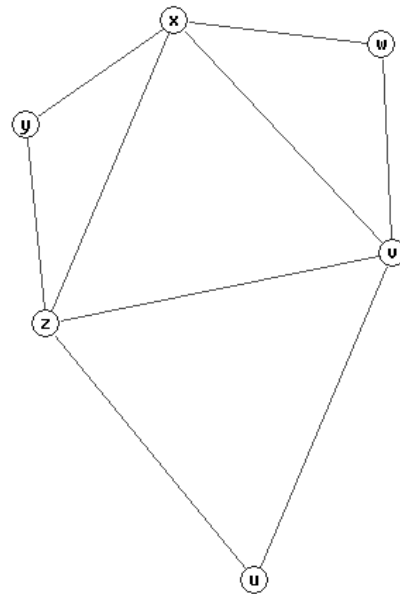
## A criterion for collapsibility of $\mathcal{M}$ onto $A$

Let

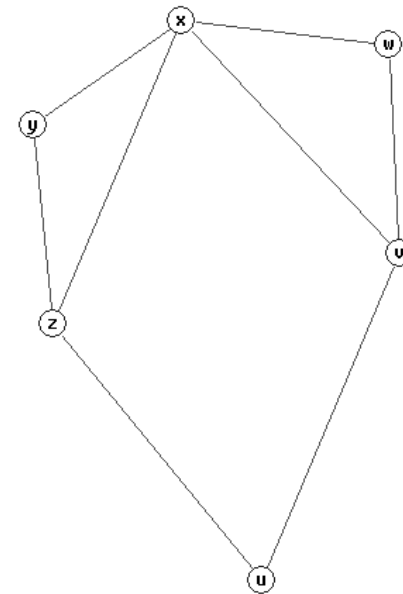
- $A \subseteq \mathcal{M}$ ,  $B = \mathcal{M} \setminus A$
- $B_1, \dots, B_k$  be the connected components of  $B$
- $E_1, \dots, E_k$  be their boundaries

Then  $\mathcal{M}$  is collapsible onto  $A$  iff each  $E_j$  is complete.

Example



Model  $\mathcal{M}$ ,  $A = \{X, Y, Z, V\}$



$\mathcal{M}$ ,  $A = \{X, Y, Z, V\}$

# Decomposability

**Definition (Lauritzen, 1996, pp 7–8)** The graphical model  $\mathcal{M}$  is *decomposable* iff

- $\mathcal{M}$  is complete; or
- $\mathcal{M}$  can be partitioned into disjoint node subsets  $A$ ,  $B$  and  $C$  such that
  - $A \perp\!\!\!\perp B|C$
  - $C$  is complete
  - $A \oplus C$  and  $B \oplus C$  are decomposable

*Note (abuse of notation):*

- $A$  can either mean the node set  $A$ , or the subgraph obtained by retaining only the edges in  $\mathcal{M}$  that begin and end with nodes in the node set  $A$ ; similarly for  $B$  and  $C$ ;
- $A \oplus B$  is the graph consisting of all the nodes in  $A$  and  $B$  and all the edges in  $\mathcal{M}$  that begin and end in  $A \cup B$ .



## Obtaining MLE's from a decomposable graph

Suppose  $\mathcal{M} = A \cup B \cup C$  where

- $A, B, C$  are disjoint
- $A \oplus C, B \oplus C$  and  $C$  are all complete
- $A \perp\!\!\!\perp B|C$

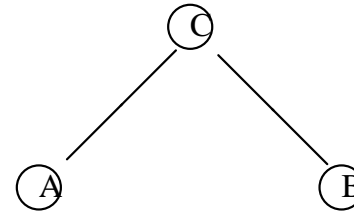
Then  $\mathcal{M}$  is decomposable, and

- $P(\mathcal{M}) = P(A, B, C) = P(A|C)P(B|C)P(C) = P(A \oplus C)P(B \oplus C)/P(C)$
- Complete graph  $\Rightarrow$  saturated model  $\Rightarrow \hat{m}_{ij\dots} = n_{ij\dots}$ , so we know how to write the MLE's for  $P(A \oplus C)$ ,  $P(B \oplus C)$  and  $P(C)$ .

If  $A \oplus C$  or  $B \oplus C$  are not complete, but are decomposable themselves, then *recurse...*

Example

$\mathcal{M} = [AC][BC]$



- $[AC]$ ,  $[BC]$ , and  $[C]$  are all complete, and  $A \perp\!\!\!\perp B|C$  by construction, so the graph is decomposable.

- The probability calculation is

$$\begin{aligned} P(\mathcal{M}) = p_{ijk} &= P(A|C)P(B|C)P(C) \\ &= P(A \oplus C)P(B \oplus C)/P(C) \\ &= p_{i+k}p_{+jk}/p_{++k} \end{aligned}$$

- Multiplying by  $m_{+++}$  on both sides, and then multiplying and dividing by  $m_{+++}$  on the right, we have

$$\begin{aligned} m_{ijk} = p_{ijk}m_{+++} &= p_{i+k}m_{+++}p_{+jk}/p_{++k} \\ &= p_{i+k}m_{+++}p_{+jk}m_{+++}/(p_{++k}m_{+++}) \\ &= m_{i+k}m_{+jk}/m_{++k} \end{aligned}$$

- For MLE's, replace  $m$ 's by  $n$ 's on the right (obs suff stats = exp suff stats).

### Interpretation

The idea of decomposability is to reduce the model formulae for  $f \in \mathcal{M}$  to a sequence of saturated models. This allows:

- Closed-form MLE's, for log-linear models (and other cases);
- Simple interpretations;
- Suggestions of a causal or time-order sequence for the variables.

Criterion [Leimer, 1989; discussed in Frydenberg & Lauritzen, 1989, Sect. 5]

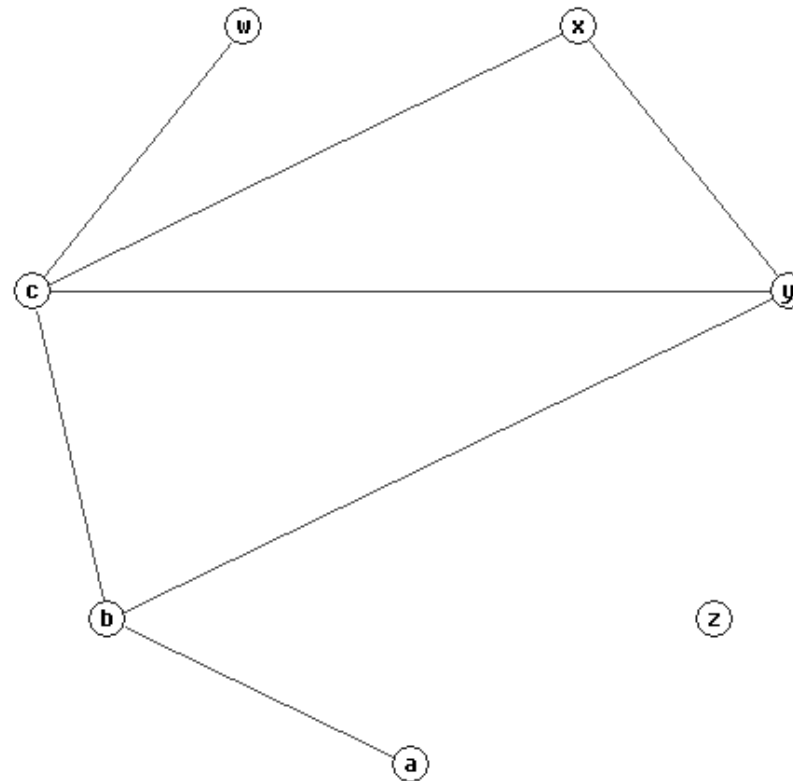
Let  $\mathcal{M}$  be a graphical log-linear model.

- Let  $c_1, \dots, c_k$  be all the cliques in  $\mathcal{M}$
- Let  $b_2 = c_2 \cap c_1, b_3 = c_3 \cap (c_2 \cup c_1), \dots, b_k = c_k \cap (\cup_{t=1}^{k-1} c_t)$  be the “running intersections sets”

The clique ordering  $c_1, \dots, c_k$  is an *SD-clique ordering* if, for every  $s = 2, 3, \dots, k$ , there is a  $j < s$  such that  $b_s \subseteq c_j$ .

The graphical model  $\mathcal{M}$  is *decomposable* iff there is an SD-clique ordering.

### Example



The cliques are  $c_1 = \{A, B\}$ ,  $c_2 = \{B, C, Y\}$ ,  $c_3 = \{C, W\}$ ,  $c_4 = \{C, X, Y\}$ ,  $c_5 = \{Z\}$ .

The running intersection sets are  $b_2 = \{B\}$ ,  $b_3 = \{C\}$ ,  $b_4 = \{C, Y\}$ ,  $b_5 = \emptyset$ .

## Some properties of decomposable graphs

- Decomposable graphs are collapsible.
  - In fact,  $\mathcal{M}$  is collapsible onto each  $u_s = \cup_{t=1}^s c_t$ ,  $s < k$ ;
  - Clearly, each  $u_s$  is also collapsible, onto each  $u_r$ ,  $r < s$ ;
- Using the relation  $f = f_{b|a}f_a$  recursively we have that

$$f = f(c_1)f(c_2|u_1)f(c_3|u_2)\cdots f(c_k|u_{k-1})$$

for all  $f \in \mathcal{M}$ .

- Thus to find the MLE  $\hat{f}$  for a decomposable log-linear graphical model we should fit the *saturated model*  $f(c_s|u_{s-1})$  to each clique, and then multiply the fitted models together.
- MLE's are not much harder to find for general decomposable models.
- It can be shown that log-linear graphical models are decomposable iff they are *triangulated*. (A stronger condition is needed if both discrete and continuous nodes are present). [Lauritzen, 1996, p. 11] See next slides.

- It turns out that decomposable graphical models are exactly those models that are “triangulated”:
  - A *cycle* is a sequence of edges  $(v_1, v_2), (v_2, v_3), \dots, (v_{k-1}, v_k)$  in  $\mathcal{M}$  such that  $v_1, \dots, v_{k-1}$  are distinct vertices, and  $v_1 = v_k$ . A *chord* is an edge  $(v_i, v_j)$  in  $\mathcal{M}$  between non-adjacent vertices ( $|i - j| > 1$ ) not already in the cycle.
  - A *triangulated graph* is one whose *chordless cycles* contain no more than three vertices.
- For a triangulated graphical model it is easy to read off the MLE’s:
  - The cell counts will have the form

$$m_{ijklpqr(etc.)} = \frac{\prod(\text{minimal sufficient margins})}{\prod(\text{separator margins})}$$

where

- \* The “*minimal sufficient margins*” are the margins with fixed indices corresponding to the terms in the generator (cliques in the graph!); and
- \* the “*separator margins*” are the margins with fixed indices corresponding to variables common to terms in the generator (these terms are *minimal* sets of vertices separating the graph into disconnected parts).

## Examples

- All log-linear models for three-way tables are decomposable, except for the model of no three-way interaction (which is not even graphical).
- [12][26][235][345]
  - Graph is triangulated
  - The separator sets are [2] (with multiplicity\* 2) and [35] (multiplicity 1)

$$m_{ijklrq} = \frac{m_{ij++++}m_{+j++++}m_{+jk+q+}m_{+++klq+}}{(m_{+j++++})^2 m_{++k+q+}}$$

- [12][26][235][345][245] not decomposable—not even graphical!
- [126][234][45][56] is graphical but not triangulated.
- [126][234][456][246] is graphical and triangulated. What are the MLE's  $\hat{m}_{ijklrq}$ ?

\*Multiplicity is one less than the number of disjoint subgraphs after removing the separator set.