

# *Network Representation Using Graph Root Distributions*

Jing Lei

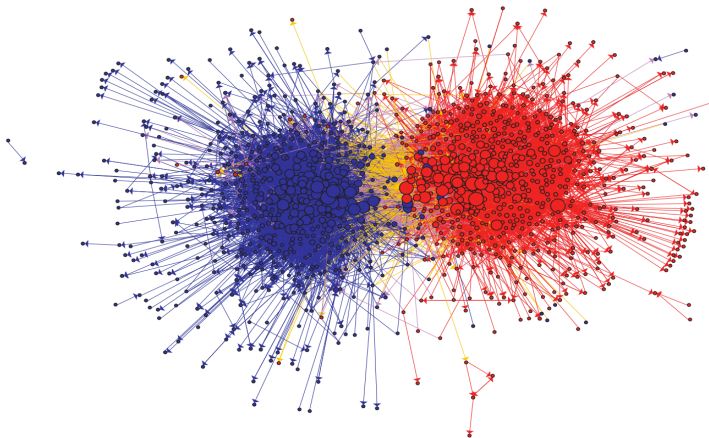
*Department of Statistics and Data Science  
Carnegie Mellon University*

2018.04

## *Network Data*

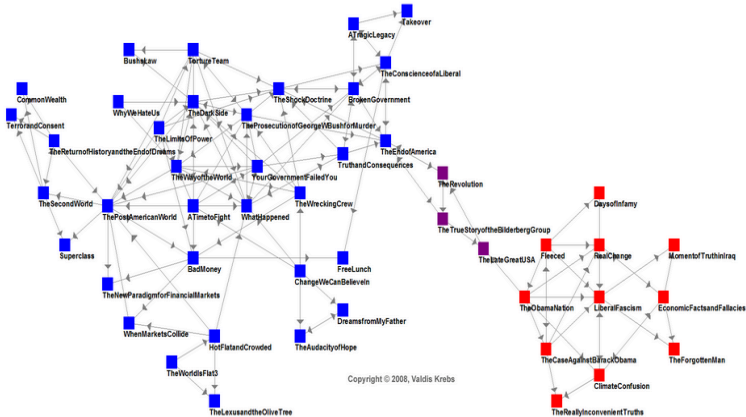
- Network data record interactions (edges) between individuals (nodes).
- From WIKIPEDIA: “... a complex network is a graph (network) with non-trivial topological features ...”
- Examples of “non-trivial topological features”
  - heavy-tail degree distribution (a.k.a “scale-free”, “power law”)
  - large clustering coefficient (transitivity)
  - community structure: the nodes can be grouped into subsets with dense internal connection.
  - ...

## *Example: Links Between Political Blogs*



[Adamic & Glance '05] The political blogosphere and the 2004 US election: divided they blog

## Example: Co-purchase of Political Books



[V. Krebs '04] Co-purchased political books on Amazon.

## Exchangeable Random Graphs

- Symmetric binary array  $\mathbf{A} = (A_{ij} : 1 \leq i < j < \infty), A_{ij} \in \{0, 1\}$ .
- Joint exchangeability:

$$(A_{ij} : 1 \leq i < j < \infty) \stackrel{d}{=} (A_{\sigma(i), \sigma(j)} : 1 \leq i < j < \infty)$$

for all permutation  $\sigma$

$$\sigma(i) = \begin{cases} i & \text{if } i \notin \{i_0, j_0\}, \\ j_0 & \text{if } i = i_0, \\ i_0 & \text{if } i = j_0. \end{cases}$$

- Idea: nodes are subjects, so the order does not matter.

## *A two-way de Finetti Theorem*

de Finetti for two-way array (Hoover '79, Aldous '81, Kallenberg '89): All such random graphs must be generated as

$$s_i \stackrel{iid}{\sim} \text{Unif}(0, 1), \quad i \geq 1.$$

$$(A_{ij}|s) \stackrel{indep.}{\sim} \text{Bernoulli}(W(s_i, s_j)), \quad 1 \leq i < j.$$

where  $W : [0, 1]^2 \mapsto [0, 1]$ , measurable and symmetric, is called a graphon (graph function).

## *Popular Special Cases*

- The stochastic block model (SBM, [Holland \*et al\* '83](#)):  $W$  is block-wise constant.
- The degree corrected block model (DCBM, [Karrer & Newman '11](#)):  $W$  is block-wise rank-one.
- Random dot product graph (RDPG, [Tang \*et al\*, '13](#); [Rubin-Delanchy \*et al\*, '17](#)):  $W$  is positive semidefinite and low-rank.
- Random geometric graphs ([Penrose '03](#)).

# *Inference Problems*

- Estimation
  - Community recovery: find block structure of  $W$  in SBM and DCBM.
  - Nonparametric estimation: estimate  $W$  from observed  $\mathbf{A}_n = (A_{ij} : 1 \leq i, j \leq n)$ .



# *Inference Problems*

- Estimation
  - Community recovery: find block structure of  $W$  in SBM and DCBM.
  - Nonparametric estimation: estimate  $W$  from observed  $\mathbf{A}_n = (A_{ij} : 1 \leq i, j \leq n)$ .
- Identifiability of  $W$ :
  - Let  $h : [0, 1] \mapsto [0, 1]$  be measure-preserving:

$$\mu(h^{-1}(B)) = \mu(B), \forall \text{ measurable } B,$$

where  $\mu$  is Lebesgue Measure.

- $W(\cdot, \cdot)$  and  $W(h(\cdot), h(\cdot))$  lead to the same distribution of  $\mathbf{A}_n$ .

## Identifiability of Graphons

- $W_1$  and  $W_2$  lead to the same distribution of  $\mathbf{A}$  if and only if there exist measure-preserving  $h_1, h_2$  such that

$$W_1(h_1(\cdot), h_1(\cdot)) = W_2(h_2(\cdot), h_2(\cdot)), \quad \text{a.e.}$$

- Cut-distance:

$$\delta_{\square}(W_1, W_2) = \inf_{h_1, h_2} \sup_{S, T} \left| \int_{S \times T} [W_1(h_1(s), h_1(t)) - W_2(h_2(s), h_2(t))] \right|$$

- When  $\delta_{\square}(W_1, W_2) = 0$ , write  $W_1 \stackrel{w.i.}{=} W_2$  (*weakly isomorphic*), which defines an equivalence relationship on  $\mathcal{W}_0 := \{W : [0, 1]^2 \mapsto [0, 1], \text{ symmetric}\}.$

## *Identifiability of Graphons*

- In general, we can only hope to recover  $W$  up to a measure-preserving change-of-variable transform.
- Existing methods assume smoothness to specify a particular member in the equivalence class (Wolfe & Olhede '13, Airoldi *et al* '13, Gao *et al* '15, Klopp *et al* '17).

## *The latent space approach*

- Sample  $\xi_1, \dots, \xi_n$  independently from a distribution  $F$  on  $\mathbb{R}^d$ .
- Connect nodes  $i, j$  by  $f(\xi_i, \xi_j)$ , for some simple function  $f$ , such as inner products and distances [Hoff et al '02, Hoff 07, Tang et al 13].
- The node embedding carries rich, interpretable structures about the network.

## *The latent space approach*

- Sample  $\xi_1, \dots, \xi_n$  independently from a distribution  $F$  on  $\mathbb{R}^d$ .
- Connect nodes  $i, j$  by  $f(\xi_i, \xi_j)$ , for some simple function  $f$ , such as inner products and distances [Hoff et al '02, Hoff 07, Tang et al 13].
- The node embedding carries rich, interpretable structures about the network.
- Question: Can we use latent space models with simple  $f$  to study exchangeable random graphs, with better identifiability?

## *The latent space approach*

- Sample  $\xi_1, \dots, \xi_n$  independently from a distribution  $F$  on  $\mathbb{R}^d$ .
- Connect nodes  $i, j$  by  $f(\xi_i, \xi_j)$ , for some simple function  $f$ , such as inner products and distances [Hoff et al '02, Hoff 07, Tang et al 13].
- The node embedding carries rich, interpretable structures about the network.
- Question: Can we use latent space models with simple  $f$  to study exchangeable random graphs, with better identifiability?
- Yes. Use graph root distributions on a separable Kreĭn space.

## Graph Root Distributions on a Kreĭn Space

### Definition: Kreĭn Space

A Kreĭn space  $\mathcal{K} = \mathcal{H}_+ \ominus \mathcal{H}_-$  is the direct sum of two Hilbert spaces  $\mathcal{H}_+$ ,  $\mathcal{H}_-$ , with inner product (for  $x, x' \in \mathcal{H}_+$ ,  $y, y' \in \mathcal{H}_-$ )

$$\langle (x; y), (x'; y') \rangle_{\mathcal{K}} = \langle x, x' \rangle_{\mathcal{H}_+} - \langle y, y' \rangle_{\mathcal{H}_-}.$$

$\mathcal{K}$  is isomorphic to a Hilbert space  $\mathcal{H} = \mathcal{H}_+ \oplus \mathcal{H}_-$  with norm  $\|\cdot\|_{\mathcal{K}}$ : for  $z = (x; y) \in \mathcal{K}$

$$\|z\|_{\mathcal{K}} = \|(x; y)\| = \left( \|x\|_{\mathcal{H}_+}^2 + \|y\|_{\mathcal{H}_-}^2 \right)^{1/2}.$$

# Graph Root Distributions on a Kreĭn Space

## Definition: Graph Root Distribution (GRD)

A graph root distribution is a probability distribution  $F$  on  $\mathcal{K}$  such that for  $Z, Z' \stackrel{iid}{\sim} F$ ,

$$\mathbb{P}(\langle Z, Z' \rangle_{\mathcal{K}} \in [0, 1]) = 1.$$



## *From GRD's to Exchangeable Random Graphs*

Given a GRD  $F$  on  $\mathcal{K}$ , one can generate exchangeable random graphs as follows.

1. Generate  $(Z_i : i \geq 1) \stackrel{iid}{\sim} F$ .
2. Generate  $A_{ij}$  independently from  $\text{Bernoulli}(\langle Z_i, Z_j \rangle_{\mathcal{K}})$ .

Related work

- Eigenmodel [[Hoff '07](#)]
- Random dot-product graph [[Tang et al '13](#); [Rubin-Delanchy et al '17](#)]

## *Interpretation of GRD*

- Edge probability =  $\langle Z_i, Z_j \rangle_{\mathcal{H}} = \langle X_i, X_j \rangle - \langle Y_i, Y_j \rangle$
- Nodes  $i, j$  are more likely to connect if
  - $\|X_i\|, \|X_j\|$  are large (active nodes)
  - $\langle X_i/\|X_i\|, X_j/\|X_j\| \rangle$  is large (good match)
- Analogous interpretations for negative components  $Y_i, Y_j$ .

## *Questions to be answered about GRD's*

- **Existence**: What kind of exchangeable random graphs can be generated by GRD's?
- **Uniqueness/Identifiability**: When do two GRD's lead to the same distribution of exchangeable random graphs?
- What is the **relationship between GRD and graphon**?
- What is an interesting **topology** in the space of GRD's?
- How to **estimate the generating GRD** from an observed network?

## *Existence of GRD Representation*

- View a graphon  $W$  as the **kernel of an integral operator** on  $L^2([0, 1])$ . By compactness,  $W$  admits spectral decomposition

$$W(s, s') = \sum_{j=1}^{\infty} \lambda_j \phi_j(s) \phi_j(s') - \sum_{j=1}^{\infty} \gamma_j \psi_j(s) \psi_j(s')$$

where  $\lambda_1 \geq \lambda_2 \geq \dots \geq 0$ ,  $\gamma_1 \geq \gamma_2 \geq \dots > 0$ .

## *Existence of GRD Representation*

### *Theorem*

If  $W$  is trace-class (i.e.  $\sum_{j \geq 1} (\lambda_j + \gamma_j) < \infty$ ), then there exists a GRD  $F$  on a separable Kreĭn space  $\mathcal{K}$  such that  $W$  and  $F$  lead to the same exchangeable random graph distribution.

## *Idea of Proof*

- Recall that

$$W(s, s') = \sum_{j=1}^{\infty} \lambda_j \phi_j(s) \phi_j(s') - \sum_{j=1}^{\infty} \gamma_j \psi_j(s) \psi_j(s')$$

- $Z(s) = (X(s); Y(s)) : [0, 1] \mapsto \mathcal{X}$  with

$$X(s) = (\lambda_j^{1/2} \phi_j(s) : j \geq 1), \quad Y(s) = (\gamma_j^{1/2} \psi_j(s) : j \geq 1).$$

- summability of  $\lambda_j, \gamma_j \Rightarrow \|X\|, \|Y\| < \infty$  a.s.  $\Rightarrow Z$  is well-defined.
- $F$  is the measure induced by  $Z$  with  $s \sim U(0, 1)$ .
- summability of  $\lambda_j, \gamma_j \Rightarrow W(\cdot, \cdot) = \langle Z(\cdot), Z(\cdot) \rangle_{\mathcal{X}}$  a.e.

## *Idea of Proof*

- Recall that

$$W(s, s') = \sum_{j=1}^{\infty} \lambda_j \phi_j(s) \phi_j(s') - \sum_{j=1}^{\infty} \gamma_j \psi_j(s) \psi_j(s')$$

- $Z(s) = (X(s); Y(s)) : [0, 1] \mapsto \mathcal{X}$  with

$$X(s) = (\lambda_j^{1/2} \phi_j(s) : j \geq 1), \quad Y(s) = (\gamma_j^{1/2} \psi_j(s) : j \geq 1).$$

- summability of  $\lambda_j, \gamma_j \Rightarrow \|X\|, \|Y\| < \infty$  a.s.  $\Rightarrow Z$  is well-defined.
- $F$  is the measure induced by  $Z$  with  $s \sim U(0, 1)$ .
- summability of  $\lambda_j, \gamma_j \Rightarrow W(\cdot, \cdot) = \langle Z(\cdot), Z(\cdot) \rangle_{\mathcal{X}}$  a.e.
- $Z$  can be viewed as the square root of  $W$ .

## *Existence of GRD Representation*

Special cases:

- Continuity:  $W = W_+ - W_-$  with continuous and positive semidefinite  $W_+, W_-$  (Mercer's theorem).
- Smoothness:  $W$  is smooth.



## Identifiability of GRD's

- When do two GRD's lead to the same sampling distribution of exchangeable random graphs?
- **Concatenation.** Let  $Z = (X; Y) \sim F$ , and  $R$  an arbitrary random variable. Let  $\tilde{Z} = (\tilde{X}; \tilde{Y})$  with  $\tilde{X} = (X, R)$ ,  $\tilde{Y} = (Y, R)$ .
- **Inner product preserving transforms.**  $H : \mathcal{H} \mapsto \mathcal{H}$  such that  $\langle z, z' \rangle_{\mathcal{H}} = \langle Hz, Hz' \rangle_{\mathcal{H}}$ . Let  $Z \sim F$  and  $\tilde{Z} = HZ$ .
  - **Direct sum of orthogonal transforms.** Let  $Q_+, Q_-$  be orthogonal transforms on  $\mathcal{H}_+, \mathcal{H}_-$ . Let  $Z = (X; Y) \sim F$ , and  $\tilde{Z} = (Q_+X; Q_-Y)$ .

## Identifiability of GRD's

- When do two GRD's lead to the same sampling distribution of exchangeable random graphs?
- **Concatenation.** Let  $Z = (X; Y) \sim F$ , and  $R$  an arbitrary random variable. Let  $\tilde{Z} = (\tilde{X}; \tilde{Y})$  with  $\tilde{X} = (X, R)$ ,  $\tilde{Y} = (Y, R)$ .
- **Inner product preserving transforms.**  $H : \mathcal{H} \mapsto \mathcal{H}$  such that  $\langle z, z' \rangle_{\mathcal{H}} = \langle Hz, Hz' \rangle_{\mathcal{H}}$ . Let  $Z \sim F$  and  $\tilde{Z} = HZ$ .
  - **Direct sum of orthogonal transforms.** Let  $Q_+, Q_-$  be orthogonal transforms on  $\mathcal{H}_+, \mathcal{H}_-$ . Let  $Z = (X; Y) \sim F$ , and  $\tilde{Z} = (Q_+X; Q_-Y)$ .
  - **Hyperbolic rotations.**

## Hyperbolic Rotations: An Example

- Let  $\mathcal{H}_+ = \mathcal{H}_- = \mathbb{R}$ ,  $z = (x; y) \in \mathbb{R}^2$ .
- An example of hyperbolic rotation is, for  $\theta \in \mathbb{R}^1$ ,

$$\begin{aligned} H[(x, y)^T] &= \left( \frac{e^\theta + e^{-\theta}}{2}x + \frac{e^\theta - e^{-\theta}}{2}y, \frac{e^\theta - e^{-\theta}}{2}x + \frac{e^\theta + e^{-\theta}}{2}y \right)^T \\ &= \begin{bmatrix} \cosh(\theta) & \sinh(\theta) \\ \sinh(\theta) & \cosh(\theta) \end{bmatrix} \begin{pmatrix} x \\ y \end{pmatrix} \\ &= H_\theta \begin{pmatrix} x \\ y \end{pmatrix}. \end{aligned}$$

- $H_\theta^T \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} H_\theta = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}.$

## *Identifiability of GRD's: Where is the Hope?*

Key observation:

- Both concatenation and hyperbolic rotation necessarily mix up the positive and negative components, so they can be precluded by disentangling the positive and negative components.

## Identifiability of GRD's

Let  $Q_+, Q_-$  be orthogonal transforms on  $\mathcal{H}_+, \mathcal{H}_-$ . Define  $Q = (Q_+ \oplus Q_-) : \mathcal{H} \mapsto \mathcal{H}$  as  $Q(x; y) = (Q_+x; Q_-y)$ .

### Theorem

Two square-integrable GRD's  $F_1, F_2$  with uncorrelated positive and negative components lead to the same sampling distribution of exchangeable random graphs if and only if there exists  $Q = Q_+ \oplus Q_-$  such that  $Z_1 \sim F_1 \Leftrightarrow Z_2 = QZ_1 \sim F_2$  (denoted as  $F_1 \stackrel{o.t.}{=} F_2$ ).

## Proof Sketch of Identifiability

- For  $i = 1, 2$ , let  $Z_i(\cdot) : [0, 1] \mapsto \mathcal{H}$  be an *inverse transform sampling (ITS)* mapping such that  $s \sim U(0, 1) \Rightarrow Z_i(s) \sim F_i$ .
- Let  $W_i(\cdot, \cdot) = \langle Z_i(\cdot), Z_i(\cdot) \rangle_{\mathcal{H}}$ .
- By assumption,  $W_1 \stackrel{w.i.}{=} W_2$ .
- Can choose appropriate orthogonal transforms so that  $W_i(\cdot, \cdot) = \langle Z_i(\cdot), Z_i(\cdot) \rangle_{\mathcal{H}}$  indeed corresponds to the spectral decomposition of  $W_i$ .
- Apply Kallenberg's representation theorem of exchangeable random arrays using spectral decompositions.

## *Existence and Identifiability of GRD: Summary*

### *Corollary*

There exists a one-to-one correspondence between trace-class graphons (up to “ $\stackrel{w.i.}{=}$ ”) and square-integrable GRD’s with uncorrelated positive and negative components (up to “ $\stackrel{o.i.}{=}$ ”).

**Canonical GRD.** Given a square integrable GRD, one can always make the positive and negative components uncorrelated and choose a canonical pair of orthogonal transforms so that the covariance is diagonalized.

## *Distances between GRD equivalence classes*

- Given two GRD's  $F_1$  and  $F_2$ , each representing their own equivalence class, how do we measure the difference between them?
- For graphons, the cut-distance is linked to the large-sample subgraph densities.

$$\delta_{\square}(W_1, W_2) = \inf_{h_1, h_2} \sup_{S, T} \left| \int_{S \times T} [W_1(h_1(s), h_1(t)) - W_2(h_2(s), h_2(t))] \right|$$

- Taking inf over  $h_1$  and  $h_2$  is to find a common ITS for two distributions, which is essentially **coupling**.



## *Wasserstein Distance*

- Let  $F_1, F_2$  be two distributions on  $\mathcal{X}$ , the Wasserstein distance between  $F_1, F_2$  is

$$d_w(F_1, F_2) = \inf_{\nu \in \mathcal{V}(F_1, F_2)} \mathbb{E}_{(Z_1, Z_2) \sim \nu} \|Z_1 - Z_2\|,$$

where  $\mathcal{V}(F_1, F_2)$  is the set of all distributions  $\nu$  on  $\mathcal{X} \times \mathcal{X}$  with marginals being  $F_1, F_2$ .

## Orthogonal Wasserstein Distance

### Definition: Orthogonal Wasserstein Distance

The *orthogonal Wasserstein distance* between two square-integrable GRD's  $F_1, F_2$  is

$$d_{\text{ow}}(F_1, F_2) = \inf_{Q_+, Q_-} \inf_{\nu \in \mathcal{V}(F_1, F_2)} \mathbb{E}_{(Z_1, Z_2) \sim \nu} \|Z_1 - (Q_+ \oplus Q_-)Z_2\|,$$

where the first inf is taken over all pairs of orthogonal transforms on  $\mathcal{H}_+, \mathcal{H}_-$ .

Remark: OWD measures the distance between two equivalence classes of GRD's.

## Stronger Topology: $d_{\text{ow}}(\cdot, \cdot) \succeq \delta_{\square}(\cdot, \cdot)$

### Theorem

Let  $F, F_1, F_2, \dots, F_N, \dots$ , be square-integrable GRD's and  $W, W_1, W_2, \dots, W_N, \dots$ , the corresponding graphons. Then

$$\delta_{\square}(W_1, W_2) \leq (\mathbb{E}_{F_1} \|Z\| + \mathbb{E}_{F_2} \|Z\|) d_{\text{ow}}(F_1, F_2).$$

Moreover, if  $d_{\text{ow}}(F_N, F) \rightarrow 0$ , then  $\delta_{\square}(W_N, W) \rightarrow 0$ .

## Estimating the GRD

- Data:  $\mathbf{A}_n = (A_{ij} : 1 \leq i, j \leq n)$ , symmetric,  $A_{ii} = 0$ .
- Model:  $A_{ij} \stackrel{\text{ind.}}{\sim} \text{Bernoulli}(\langle Z_i, Z_j \rangle_{\mathcal{H}})$ ,  $1 \leq i < j \leq n$ , where  $(Z_i : 1 \leq i \leq n) \stackrel{\text{iid}}{\sim} F$ , a square-integrable GRD on a  $\mathcal{H}$ .
- Let  $\mathcal{H}_+ = \mathcal{H}_- = \{x \in \mathbb{R}^{\mathbb{N}} : \sum_{j \geq 1} x_j^2 < \infty\}$ .
- To identify, let  $Z = (X; Y) \sim F$  have diagonal covariance:

$$\mathbb{E}XX^T = \text{diag}(\lambda_1, \lambda_2, \dots), \quad \mathbb{E}YY^T = \text{diag}(\gamma_1, \gamma_2, \dots), \quad \mathbb{E}XY^T = 0.$$

## Truncated Weighted Signed Spectral Embedding

- Let  $\mathbf{A}_n = \sum_j \hat{\lambda}_j \hat{a}_j \hat{a}_j^T - \sum_j \hat{\gamma}_j \hat{b}_j \hat{b}_j^T$  be the eigen decomposition of  $\mathbf{A}_n$ , with positive eigenvalues  $\hat{\lambda}_j$  and (absolute) negative eigenvalues  $\hat{\gamma}_j$  ranked in decreasing order.
- Let  $p_1, p_2$  be positive integers to be specified later.
- For  $1 \leq i \leq n$  let  $\hat{Z}_i = (\hat{X}_i; \hat{Y}_i)$  with
$$\hat{X}_i = (\hat{\lambda}_1^{1/2} \hat{a}_{1i}, \dots, \hat{\lambda}_{p_1}^{1/2} \hat{a}_{p_1 i}, 0, \dots)$$
$$\hat{Y}_i = (\hat{\gamma}_1^{1/2} \hat{b}_{1i}, \dots, \hat{\gamma}_{p_2}^{1/2} \hat{b}_{p_2 i}, 0, \dots)$$
- $\hat{F}$  is the distribution that puts  $1/n$  mass at each  $\hat{Z}_i$ .

## Regularity Conditions

- Eigen decay and eigen gap: for some  $1 < \alpha \leq \beta$  and all  $j \geq 1$

$$\lambda_j, \gamma_j \asymp j^{-\alpha}, \quad \min(\lambda_j - \lambda_{j+1}, \gamma_j - \gamma_{j+1}) \gtrsim j^{-\beta}$$

- Fourth moment bounded:  $\mathbb{E}_{Z \sim F} \|Z\|^4 < \infty$ .
- These are standard assumptions in functional data analysis, where truncated PCA is used to recover sample curves in  $L^2([0, 1])$ .

## Estimation Error Bound

Let  $\tilde{F}$  be the distribution that puts  $1/n$  mass at  $\tilde{Z}_i = (\tilde{X}_i, \tilde{Y}_i)$ , with  $\tilde{X}_i = (X_{i1}, \dots, X_{ip}, 0, \dots)$ ,  $\tilde{Y}_i = (Y_{i1}, \dots, Y_{ip}, 0, \dots)$

### Theorem

Under the regularity conditions, if  $p_1 = p_2 = p = o(n^{1/(2\beta+\alpha)})$  then

$$d_w(\hat{F}, \tilde{F}) = O_P(p^{-(\alpha-1)/2})$$

and

$$d_w(\hat{F}, F) = O_P(p^{-(\alpha-1)/2} + n^{-1/p}).$$

## Proof Sketch

- Treat positive and negative components separately.
- $X_i = (X_{ij} : j \geq 1)$   
 $\hat{X}_i = (\hat{\lambda}_1^{1/2} \hat{a}_{1i}, \dots, \hat{\lambda}_p^{1/2} \hat{a}_{pi}, 0, \dots)$   
 $\tilde{X}_i = (X_{i1}, \dots, X_{ip}, 0, \dots)$
- $\mathbf{G}_n = (\langle (X_i; Y_i), (X_j; Y_j) \rangle_{\mathcal{H}} : 1 \leq i, j \leq n) = \mathbb{E}(\mathbf{A}_n | Z_1, \dots, Z_n).$
- $\mathbf{G}_{n,X} = (\langle X_i, X_j \rangle : 1 \leq i, j \leq n), \mathbf{G}_n = \mathbf{G}_{n,X} - \mathbf{G}_{n,Y}$
- $\tilde{\mathbf{G}}_{n,X} = (\langle \tilde{X}_i, \tilde{X}_j \rangle : 1 \leq i, j \leq n)$
- Spectral perturbation:  $\mathbf{A}_n \approx \mathbf{G}_n$   
 $\Rightarrow \hat{X}_i \approx \text{T.W. spectral embedding of } \mathbf{G}_n.$
- Uncorrelatedness+eigen-decay:  $\mathbf{G}_n \approx \mathbf{G}_{n,X} \approx \tilde{\mathbf{G}}_{n,X}$  in leading subspace  $\Rightarrow \hat{X}_i \approx \tilde{X}_i \Rightarrow d_w(\hat{F}, \tilde{F}) \approx 0.$



## Sparse Graphs

- GRD's, as graphons, can only generate dense graphs with total number of edges proportional to  $n^2$ .
- Given graphon  $W$  and node sample size  $n$ , one can consider sparse sampling with a sparsity rate  $\rho_n = o(1)$  (see e.g. [\[Bickel & Chen '09\]](#)):

$$A_{ij} \sim \text{Bernoulli}(\rho_n W(s_i, s_j)).$$

- In GRD representation, the sparse sampling is equivalent to generating the network using  $\rho_n^{1/2} F$  (scaling down  $F$  by a factor of  $\rho_n^{1/2}$ ).

# Sparse Graphs

## Theorem

Assume  $\mathbf{A}_n$  is generated by a GRD  $F$  with sparsity parameter  $\rho_n$ . Under the regularity conditions, if  $\beta \geq 3\alpha/2$  and  $p = o(n^{1/(2\beta+\alpha)} \wedge (n\rho_n)^{1/(2\beta)})$  then

$$d_w(\rho_n^{-1/2}\hat{F}, F) = O_P\left(\frac{p^{\beta-(\alpha-1)/2}}{(n\rho_n)^{1/2}} + p^{-(\alpha-1)/2} + n^{-1/p}\right).$$

Other values of  $\beta$  and  $p$  are allowed, but complicated to present.

## *How to Choose $p_1, p_2$ ?*

- The truncated empirical eigen decomposition resembles methods in functional principal components analysis, where one can choose the number of PC's by fraction of variance explained.
- However, network data are different
  - Low-rank: the number of significant eigen components is usually small;
  - High noise: network data are observed with entry-wise independent noise.
- Singular value thresholding [[Chatterjee '14](#)]: use eigenvalues larger than  $\sqrt{n}$ .

## Examples

- $B \in [0, 1]^{K \times K}$ ,  $B = B^T$ .
- Stochastic block models: mixture of point mass

$$\mathbb{E}(A_{ij}) = B_{g_i, g_j}, \quad g_i \in \{1, \dots, K\}.$$

- Degree corrected block models: mixture of 1-D subspaces

$$\mathbb{E}(A_{ij}) = \psi_i \psi_j B_{g_i, g_j}, \quad g_i \in \{1, \dots, K\}, \quad \psi_i \in \mathbb{R}^+.$$

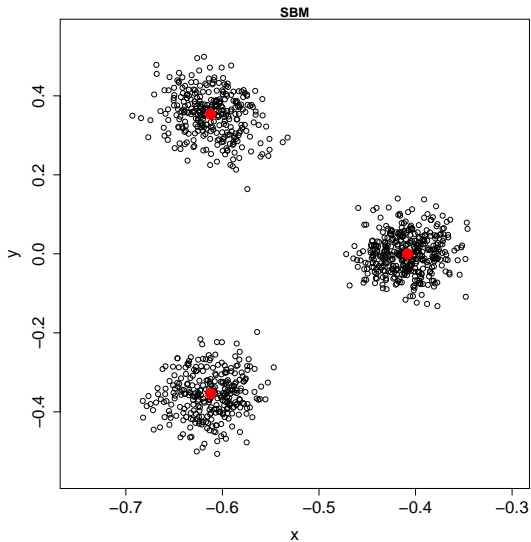
- Mixed membership block models: convex polytope

$$\mathbb{E}(A_{ij}) = a_i^T B a_j, \quad a_i \in (K-1) \text{ dim. simplex}.$$

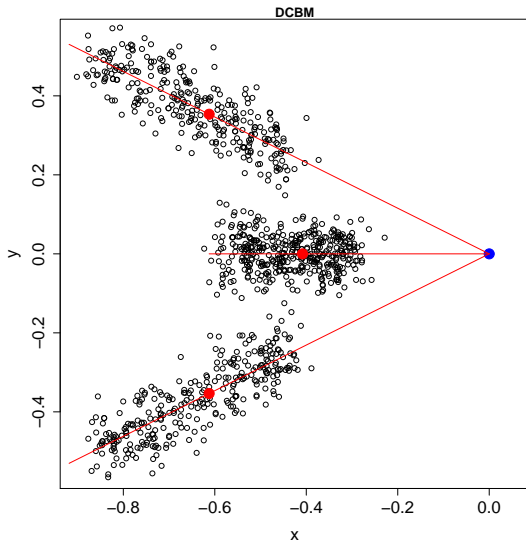
## Simulation, $K = 3$

- $K = 3$ ,  $g_i \sim \text{Multinomial}(0.3, 0.3, 0.4)$ ,  $n = 1000$ .
- $B = \begin{bmatrix} 1/4 & 1/2 & 1/4 \\ 1/2 & 1/4 & 1/4 \\ 1/4 & 1/4 & 1/6 \end{bmatrix}$
- Three communities but  $\text{rank}(B) = 2$ , with one positive and one negative component.
- DCBM node activeness:  $\psi_i \sim U(0.7, 1.4)$ .
- MMBM node mixture:  $a_i \sim \text{Dir}(0.5, 0.5, 0.5)$ .

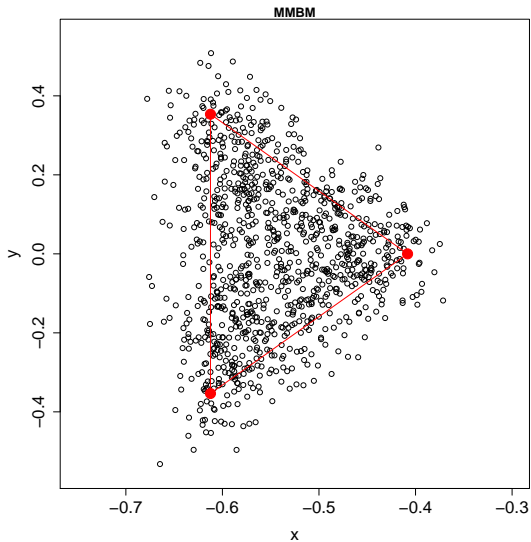
## *SBM: Point Mass Mixture*



## *DCBM: Subspace Mixture*

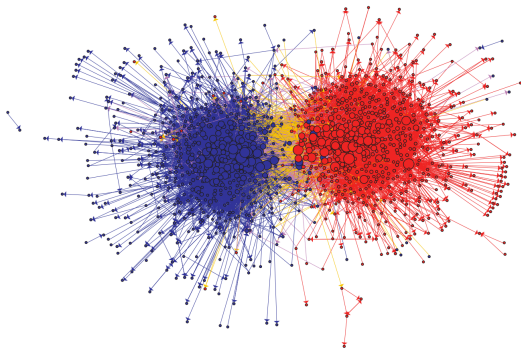


## *MMBM: Convex Polytope*



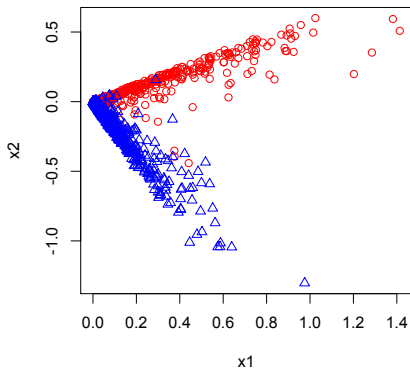
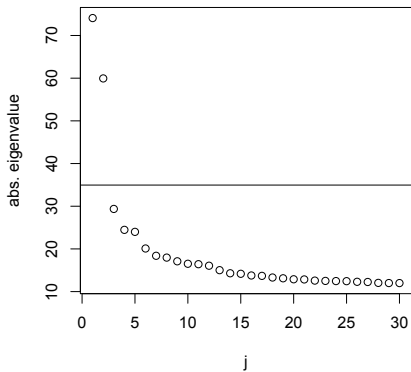


## *Data example: U.S. political blogs*

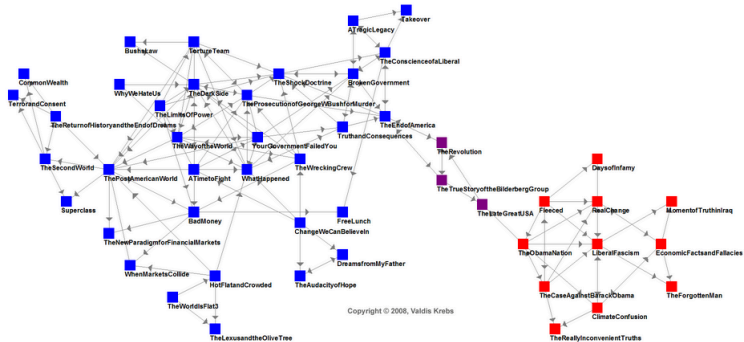


- [Adamic & Glance '05] Snapshots of weblogs shortly before 2004 U.S. Presidential Election. Nodes: weblogs; edges: hyperlinks.
- Fitted well by a DCBM with two clusters.

## *Political Blogs Data*

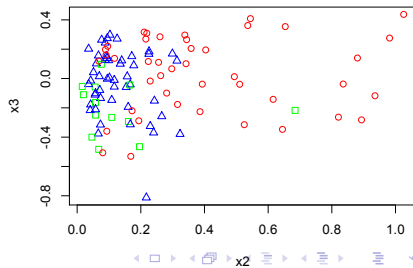
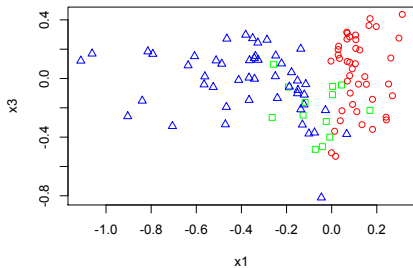
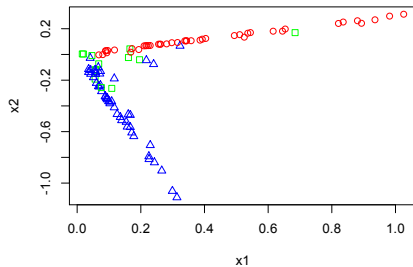
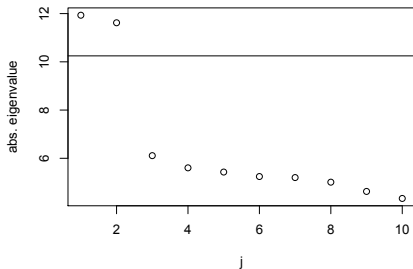


# Political Books Data



Co-purchase of political books on Amazon (V. Krebs '04)

# Political Books Data



## Next

- GRD with logistic link:  $A_{ij} \sim \text{Bernoulli} \left( \frac{1}{1 + e^{-(Z_i, Z_j)_{\mathcal{K}}}} \right)$
- Two sample testing: are  $\mathbf{A}_1$  and  $\mathbf{A}_2$  from the same GRD?
- Bi-partite graph:  $\mathcal{A}$  is asymmetric. e.g. gene-cell matrix.
- Multiple networks: modeling and predicting the movement of latent node embeddings.

## *Thank You! Questions?*

1. Lei, J. “Network Representation Using Graph Root Distributions”, `arXiv:1802.09684`
2. Code: easy to write but also available upon request
3. Slides: `www.stat.cmu.edu/~jinglei/talk.shtml`