

Administrivia

Simplex Algorithm in 1 Slide

Canonical form:

$$\begin{array}{l}
 -Z \\
 x_1 \quad x_2 \quad \vdots \quad x_m \\
 + C_{m+1}x_{m+1} + \dots + C_jx_j + \dots + C_nx_n = -Z_0 \\
 + a_{1,m+1}x_{m+1} + \dots + a_{1j}x_j + \dots + a_{1n}x_n = b_1 \\
 \vdots \\
 x_m + a_{m,m+1}x_{m+1} + \dots + a_{mj}x_j + \dots + a_{mn}x_n = b_m
 \end{array}$$

If we do pivot in $A_{r,s} > 0$, where $c_s < 0$

$$\boxed{X_s \text{ IN } \quad X_r \text{ OUT}}$$

New cost value:

$$\tilde{Z} = Z_0 + \underbrace{b_r}_{\substack{0 \\ \nearrow}} \underbrace{\frac{C_s}{A_{rs}}}_{\substack{0 \\ \downarrow \\ 0}}$$

CONSTRAINTS

$$\begin{aligned}
 x_1 &= b_1 - A_{1s}x_s \geq 0 \\
 &\vdots \\
 x_m &= b_m - A_{ms}x_s \geq 0 \\
 \Rightarrow x_s &= \min \left(\frac{b_i}{A_{is}} \right)
 \end{aligned}$$

New b vector:

$$\tilde{b}_j = b_j + b_r \cdot \frac{-A_{js}}{A_{rs}} \geq 0 \Rightarrow \text{LET } r = \text{ARG MIN}_{\{i | A_{is} > 0\}} \frac{b_i}{A_{is}}$$

The full Simplex Algorithm

So far we have assumed that a basic feasible solution in canonical form is available to start the algorithm...

The Simplex Algorithm Phase I

Example

$$\begin{array}{l}
 2x_1 + 1x_2 + 2x_3 + x_4 + 4x_5 = \infty \rightarrow \left. \begin{array}{l} \min \\ x_i \geq 0 \end{array} \right\} (*1) \\
 \text{s.t. } 4x_1 + 2x_2 + 13x_3 + 3x_4 + x_5 = 17 \\
 x_1 + x_2 + 5x_3 + x_4 + x_5 = 7
 \end{array}$$

Goal: We want to find a feasible solution

Phase I:

- (i) Forget the cost function $c^T x$.
- (ii) Introduce $x_6, x_7 \geq 0$. [One variable for each row]

$$\begin{array}{l}
 \text{(iii) Solve } \min x_6 + x_7 = W \\
 x_i \geq 0 \quad i=1..7 \\
 \text{s.t. } 4x_1 + 2x_2 + 13x_3 + 3x_4 + x_5 + x_6 + 0x_7 = 17 \\
 x_1 + x_2 + 5x_3 + x_4 + x_5 + 0x_6 + x_7 = 7
 \end{array}$$

The Simplex Algorithm Phase I

$$\begin{array}{l}
 \text{MIN } X_6 + X_7 = W \\
 X_i \geq 0 \quad i=1..7 \\
 \text{s.t. } \left. \begin{array}{l}
 4X_1 + 2X_2 + 13X_3 + 3X_4 + X_5 + X_6 + 0X_7 = 17 \\
 X_1 + X_2 + 5X_3 + X_4 + X_5 + 0X_6 + X_7 = 7
 \end{array} \right\} (*2)
 \end{array}$$

Theorem: (*2) has feasible optimal solution such that $X_6=X_7=0$
iff (*1) has feasible solution

Remarks:

- (*2) is easy to convert to a feasible canonical solution (We will see)
- We can find its optimal solution ($X_6=X_7=0$) with the Phase II algorithm
This is a feasible solution of (*1)

Tableaux

$$\begin{array}{l}
 \text{MIN } X_6 + X_7 = W \\
 x_i \geq 0 \quad i=1..7 \\
 \text{s.t. } \left. \begin{array}{l} 4x_1 + 2x_2 + 13x_3 + 3x_4 + x_5 + x_6 + 0x_7 = 17 \\ x_1 + x_2 + 5x_3 + x_4 + x_5 + 0x_6 + x_7 = 7 \end{array} \right\} (\times 2)
 \end{array}$$

Basic variable	Objective (-w)	x_1	x_2	x_3	x_4	x_5	x_6	x_7	RHS
	1	0	0	0	0	0	1	1	0
	0	4	2	13	3	1	1	0	17
	0	1	1	5	1	1	0	1	7

g-

It's easy to convert this to a feasible canonical form:

Basic variable	Objective (-w)	x_1	x_2	x_3	x_4	x_5	x_6	x_7	RHS
-w	1	-5	-3	-18	-4	-2	0	0	-24
x_6	0	4	2	13	3	1	1	0	17
x_7	0	1	1	5	1	1	0	1	7

CANONICAL FORM

Tableaux

Now, we can run the Phase 2 algorithm on this table to get a feasible solution of (*1):

B. var	Obj. (-w)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-w	1	-5	-3	-18	-4	-2	0	0	-24
X_6	0	4	2	13	3	1	1	0	17
X_7	0	1	1	5	1	1	0	1	7

-18 is the smallest $C_j < 0 \Rightarrow$ pivot in column 3

$\frac{17}{13} < \frac{7}{5} \Rightarrow$ pivot in 13
 X_3 IN
 X_6 OUT

Tableaux

Now, we can run the Phase 2 algorithm on this table to get a feasible solution of (*1):

B. var	Obj. (-w)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-w	1	-5	-3	-18	-4	-2	0	0	-24
X_6	0	4	2	13	3	1	1	0	17
X_7	0	1	1	5	1	1	0	1	7



B. var	Obj. (-w)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-w	1	$-5 + 18/13 \cdot 4$	$-3 + 18/13 \cdot 2$	0	$-4 + 18/13 \cdot 3$	$-2 + 18/13$	$18/13$	0	$-24 + 18/13 \cdot 7$
X_3	0	$4/13$	$2/13$	1	$3/13$	$1/13$	$1/13$	0	$17/13$
X_7	0	$1 - 5/13 \cdot 4$	$1 - 5/13 \cdot 2$	0	$1 - 5/13 \cdot 3$	$1 - 5/13$	$-5/13$	1	$7 - 5/13 \cdot 7$

Tableaux

B. var	Obj (-w)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-w	1	$-5+18/13*4$	$-3+18/13*2$	0	$-4+18/13*3$	$-2+18/13$	$18/13$	0	$-24+18/13*7$
X_3	0	$4/13$	$2/13$	1	$3/13$	$1/13$	$1/13$	0	$17/13$
X_7	0	$1-5/13*4$	$1-5/13*2$	0	$1-5/13*3$	$1-5/13$	$-5/13$	1	$7-5/13*7$

Let us simplify this Table a little:

B. var	Obj (-w)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-w	1	$-7/13$	$-3/13$	0	$2/13$	$-8/13$	$18/13$	0	$-6/13$
X_3	0	$4/13$	$2/13$	1	$3/13$	$1/13$	$1/13$	0	$17/13$
X_7	0	$1-7/13$	$3/13$	0	$-2/13$	$8/13$	$-5/13$	1	$6/13$

↑
PIVOT

$-8/13$ IS THE SMALLEST AMONG ALL $C_j < 0 \Rightarrow X_5$ IN

$$\frac{6/13}{8/13} < \frac{17/13}{1/13} \Rightarrow X_7 \text{ OUT}$$

Tableaux

B. var	Obj (-w)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-w	1	0	0	0	0	0	1	1	0
X_3	0	3/8	1/8	1	1/4	0	-1/8	0	5/4
X_5	0	-7/8	3/8	0	-1/4	1	-5/8	1/8	3/4

All the relative costs are nonnegative \Rightarrow optimal feasible solution.

Phase I is finished.

$-w + x_6 + x_7 = 0$
 $x_3 = 5/4$
 $x_5 = 3/4$
 $x_1, x_2, x_4, x_6, x_7 = 0$

FEASIBLE BASIC SOLUTION
 OF THE ORIGINAL (*) PROBLEM
 \Downarrow
 PHASE II CAN BE STARTED

Tableaux

B. var	Obj (-w)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-w	1	0	0	0	0	0	1	1	0
X_3	0	3/8	1/8	1	1/4	0	-1/8	0	5/4
X_5	0	-7/8	3/8	0	-1/4	1	-5/8	1/8	3/4

ORIGINAL COST FUNCTION: $z = 2x_1 + 1x_2 + 2x_3 + x_4 + 4x_5 \rightarrow \min_{x \geq 0}$

B. var	Obj (-z)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-z	1	2	1	2	1	4	0	0	0
X_3	0	3/8	1/8	1	1/4	0	-1/8	0	5/4
X_5	0	-7/8	3/8	0	-1/4	1	-5/8	1/8	3/4

Tableaux

B. var	Obj (-z)	X_1	X_2	X_3	X_4	X_5	X_6	X_7	RHS
-z	1	2	1	2	1	4	0	0	0
X_3	0	$3/8$	$1/8$	1	$1/4$	0	$-1/8$	0	$5/4$
X_5	0	$-7/8$	$3/8$	0	$-1/4$	1	$-5/8$	$1/8$	$3/4$

Make it to canonical form and continue with Phase II...

Do not make pivots in the column of X_6 and X_7

Simplex Algorithm with Matlab

```
f = [-5 -4 -6]';
```

```
A = [ 1 -1 1  
      3 2 4  
      3 2 0];
```

```
b = [20 42 30]';
```

```
lb = zeros(3,1);
```

```
options = optimset('LargeScale','off','Simplex','on');
```

```
[x,fval,exitflag,output,lambda] = linprog(f,A,b,[],[],lb,[],[],options);
```

Relevant Books

- ❑ Luenberger, David G. *Linear and Nonlinear Programming*. 2nd ed. Reading, MA: Addison Wesley, 1984. ISBN: 0201157942.
- ❑ Bertsimas, Dimitris, and John Tsitsiklis. *Introduction to Linear Optimization*. Athena Scientific Press, 1997. ISBN: 1886529191.
- ❑ Dantzig, Thapa: Linear Programming

Summary

- ❑ Linear programs:
 - standard form,
 - canonical form
- ❑ Solutions:
 - Basic, Feasible, Optimal, Degenerate
- ❑ Simplex algorithm:
 - Phase I
 - Phase II
- ❑ Applications:
 - Pattern classification

Convex Optimization

CMU-10725

4. Convexity Part I

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MACHINE LEARNING DEPARTMENT



Goal of this lecture

□ Review of Convex sets & Convex functions

- Definition
- Examples
- Basic properties

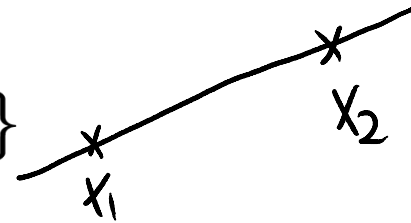
Books to Read:

- **Boyd and Vandenberghe**: Convex Optimization, Chapters 2 & 3
- **Rockafellar**: Convex Analysis

Line and Line Segments

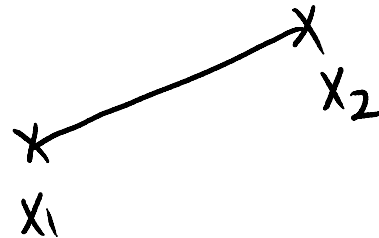
Definition [Line]:

$$\{x \in \mathbb{R}^n | x = \theta x_1 + (1 - \theta)x_2, \theta \in \mathbb{R}\}$$



Definition [Line segment]: $\text{LINE} + \theta \in [0, 1]!$

$$\{x \in \mathbb{R}^n | x = \theta x_1 + (1 - \theta)x_2, \theta \in [0, 1]\}$$



Affine Sets

Definition [Affine set]:

A set C is affine if for any $x_1, x_2 \in C$
the line through x_1 and x_2 is in C ,
i.e. $\theta x_1 + (1 - \theta)x_2 \in C$, ($\theta \in \mathbb{R}$)

Definition [Affine hull of set C]:

$$C \subseteq \mathbb{R}^n \quad \text{AFF}[C] = \left\{ e_1 x_1 + \dots + e_n x_n \mid x_1, \dots, x_n \in C, \sum_{i=1}^n e_i = 1 \right\}$$

Theorem [Affine hull]:

The $\text{Aff}[C]$ is the smallest affine set that contains C

Affine Sets Example

Example [Solutions of linear equations]:

The solution set of a system of linear equations is an affine set

Solution set:

$$C = \{x \mid Ax = b\} \quad A \in \mathbb{R}^{m \times n}, \quad b \in \mathbb{R}^m$$

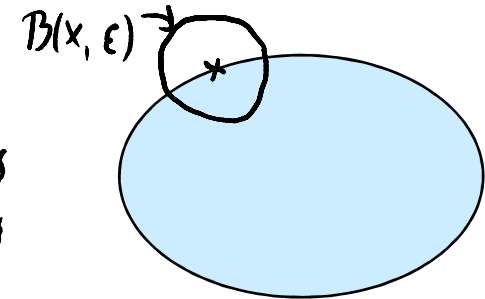
Proof:

$$\left. \begin{array}{l} \text{If } Ax_1 = b \\ Ax_2 = b \\ \alpha \in \mathbb{R} \end{array} \right\} \Rightarrow A[\alpha x_1 + (1-\alpha)x_2] = b$$

Boundaries

Definition [x on boundary of C (∂C)]:

FOR SMALL ENOUGH $\epsilon > 0$: $B(x, \epsilon) \cap C \neq \emptyset$
 & $B(x, \epsilon) \cap C^c \neq \emptyset$

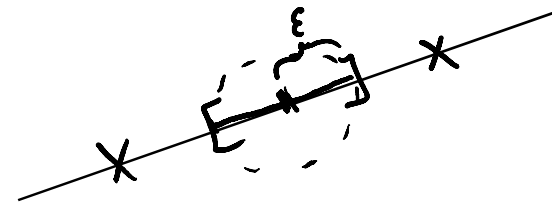


Definition [x in interior of C]:

$B(x, \epsilon) \subset C$ FOR SMALL ENOUGH ϵ

Definition [*relative* interior (rel int C)]:

$\text{REL INT } C = \{ x \in C \mid B(x, \epsilon) \cap \text{AFF } C \subseteq C \text{ FOR SOME } \epsilon > 0 \}$



Boundaries

Definition [closure of C ($\text{cl } C$)]:

$$\text{cl } C = C \cup \partial C$$

Definition [*relative* boundary of C ($\text{rel } \partial C$)]:

$$\text{cl } C \setminus \text{REL INT } C$$

Open and Closed Sets

Definition [C closed]:

$$\partial C \subset C$$

Definition [C open]:

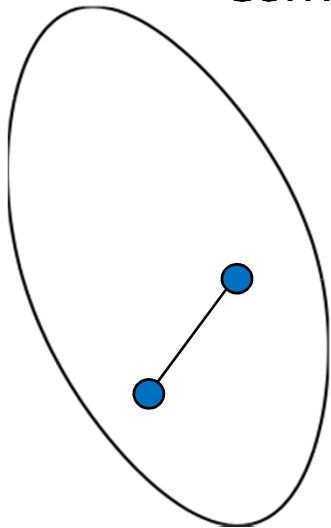
$$\partial C \cap C = \emptyset$$

Definition [C compact]:

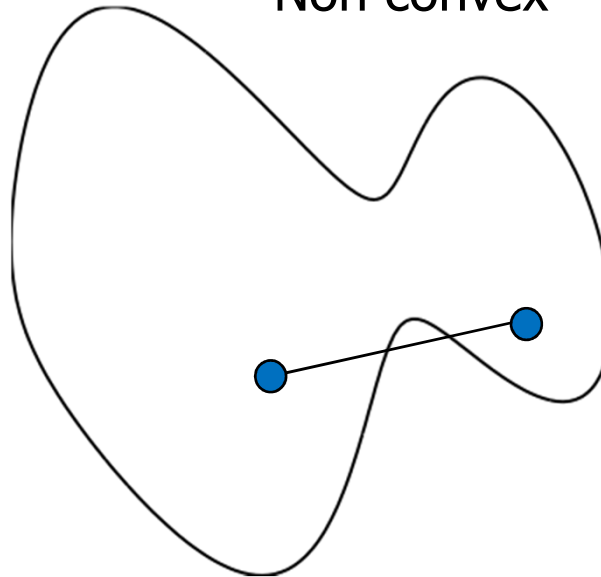
CLOSED AND BOUNDED [IN \mathbb{R}^n]

Convex sets

Convex



Non convex



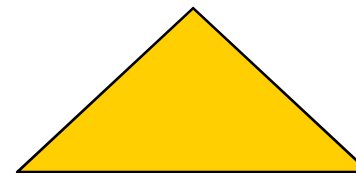
Definition [Convex set]:

SET $C \subset \mathbb{R}^d$ is convex $(\Leftrightarrow) \forall x_1, x_2 \in C \forall \theta \in [0, 1]$
 $\theta x_1 + (1 - \theta)x_2 \in C$

Definition [Strictly convex set]:

$\theta x_1 + (1 - \theta)x_2 \in \text{INT}(C) \forall x_1 \neq x_2 \in C \theta \in (0, 1)$

Example [Convex, but not strictly convex]:



Examples of Convex Sets

- empty set: \emptyset
- singleton set: $\{x_0\}$
- complete space: \mathbb{R}^d
- lines: $\{x \mid x = \theta x_1 + (1-\theta)x_2, \theta \in \mathbb{R}\}$
- line segments: $\{x \mid x = \theta x_1 + (1-\theta)x_2, \theta \in [0, 1]\}$
- hyperplanes: $\{x \in \mathbb{R}^n \mid a^T x = b\} \quad a \in \mathbb{R}^n, b \in \mathbb{R}$
- halfspaces: $\{x \in \mathbb{R}^n \mid a^T x \leq b\} \quad a \in \mathbb{R}^n, b \in \mathbb{R}$

Examples of Convex Sets

- Euclidian balls:

$$\{x \in \mathbb{R}^n \mid \|x - x_0\|_2 \leq r\} = \mathcal{B}(x_0, r)$$

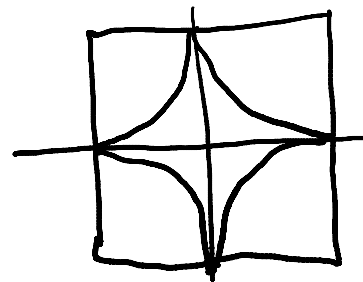
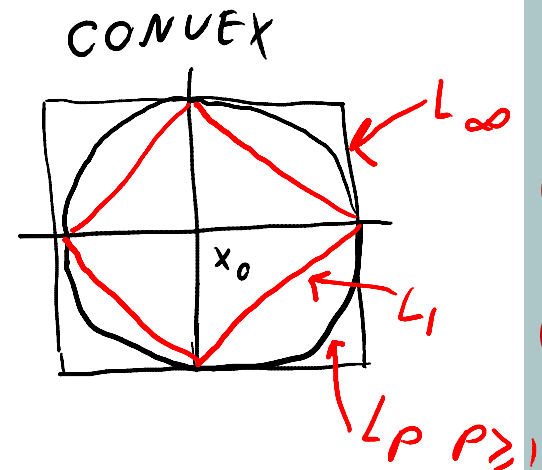
- L_p balls, $p \geq 1$

$$\{x \in \mathbb{R}^n \mid \|x - x_0\|_p \leq r\}$$

$$\|x\|_p = \left(\sum_{i=1}^n |x_i|^p \right)^{1/p}$$

$$\|x\|_\infty = \max_{i=1..n} |x_i|$$

- L_p balls $0 < p < 1$



NOT CONVEX!

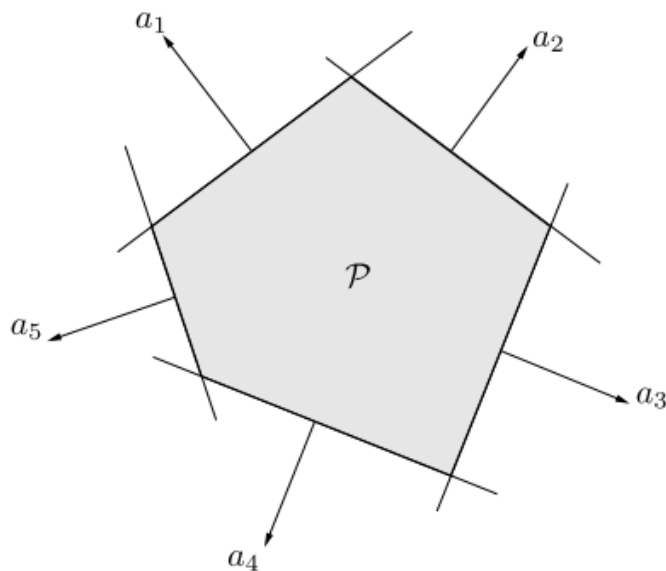
Examples of Convex Sets

- Polyhedron: the solution set of a finite number of linear equalities and inequalities

$$\mathcal{P} = \{x \mid a_j^T x \leq b_j, j=1, \dots, m \\ c_j^T x = d_j, j=1, \dots, p\}$$

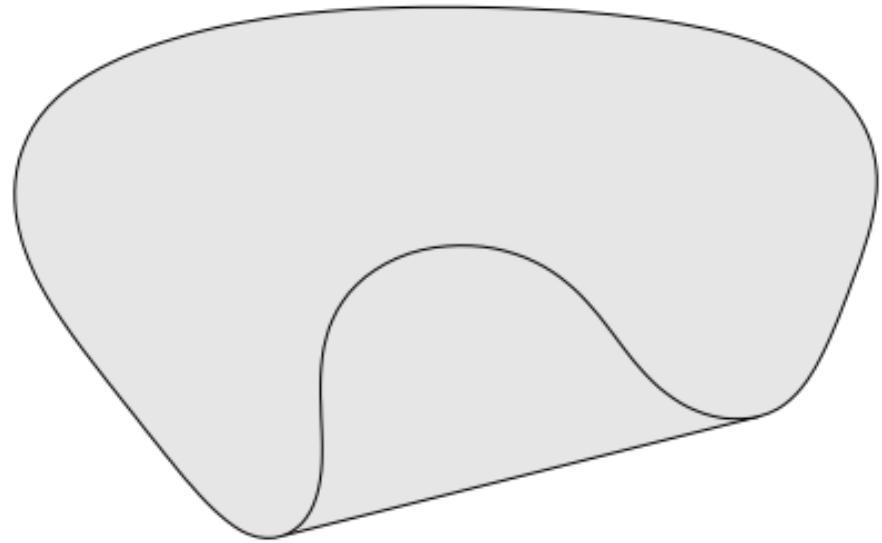
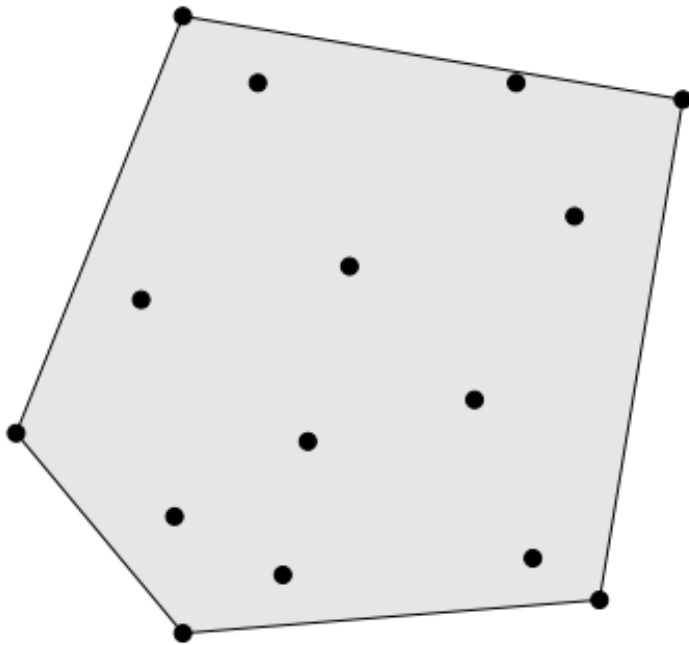
Matrix notation: $\mathcal{P} = \{x \mid Ax \leq b, Cx = d\}$

- Polytope: bounded polyhedron



Intersection of halfspaces
& hyperplanes

Convex hull



Definition [Convex hull]:

$$\text{conv}[C] = \left\{ \theta_1 x_1 + \dots + \theta_k x_k \mid x_i \in C, \theta_i \geq 0 \ \forall i=1 \dots k, \sum_{i=1}^k \theta_i = 1, k \in \mathbb{Z}_+ \right\}$$

Convex hull

Convex hull properties:

$\text{CONV}[C]$ IS THE SMALLEST CONVEX SET
THAT CONTAINS C

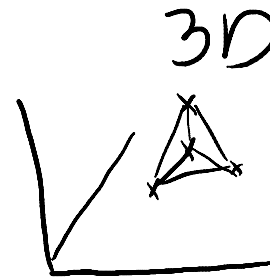
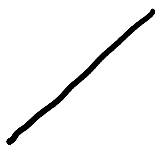
$\text{CONV}[C]$ IS CONVEX

$C \subseteq \text{CONV}(C)$

$\forall C, C'$ CONVEX SETS. IF $C \subset C' \Rightarrow \text{CONV}[C] \subseteq C'$

Examples of Convex Sets

□ Simplex:



Convex Combination

$$\text{CONV}[C] = \left\{ \theta_1 x_1 + \dots + \theta_k x_k \mid x_i \in C, \theta_i \geq 0 \forall i=1 \dots k, \sum_{i=1}^k \theta_i = 1, k \in \mathbb{Z}_+ \right\}$$

Infinite many sums

$$\begin{aligned} C \text{ CONVEX} &\Rightarrow \underbrace{\sum_{i=1}^{\infty} \theta_i x_i}_{\text{ASSUMING THIS SERIES CONVERGES}} \in C \\ x_1, x_2, \dots &\in C \\ \theta_1, \theta_2, \dots &\geq 0 \\ \sum_{i=1}^{\infty} \theta_i &= 1 \end{aligned}$$

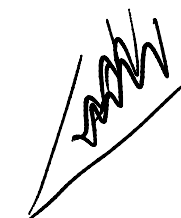
Integrals (Expected value)

$$\begin{aligned} P: \mathbb{R}^n &\rightarrow \mathbb{R} \text{ DENSITY FUNCTION} \\ C \text{ CONVEX SET} &\Rightarrow \underbrace{\int_C P(x) x dx}_{\text{ASSUMING THIS INTEGRAL EXISTS}} \in C \\ \mathbb{E}_P[X] &\in C \end{aligned}$$

Cones

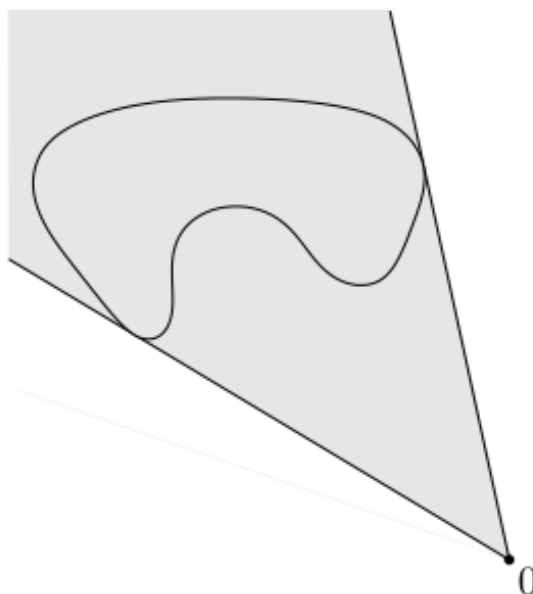
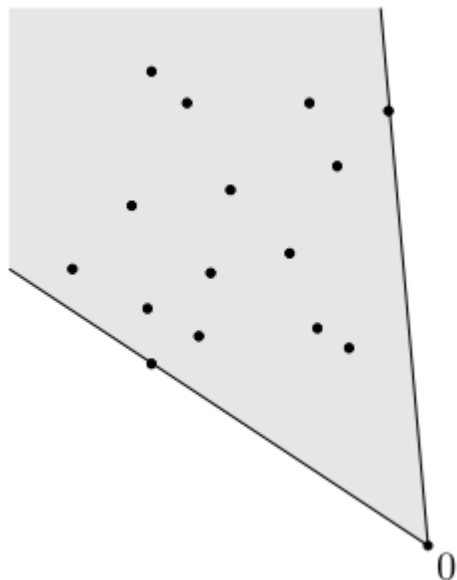
Definition [Cone]:

$$C \text{ is CONE} \Leftrightarrow \left\{ \begin{array}{l} x \in C \\ \theta \geq 0 \end{array} \Rightarrow \theta x \in C \right\}$$



Definition [Convex Cone]: CONE & CONVEX

Definition [Conic hull]: $\text{CONE}[C] = \left\{ x \mid x = \theta_1 x_1 + \dots + \theta_k x_k, \theta_i \geq 0, x_i \in C \right\}$



Example: PSD matrices

Definition [Positive semi definite matrix]:

$$A \in \mathbb{R}^{n \times n} \text{ MATRIX IS PSD } \Leftrightarrow \begin{cases} x^T A x \geq 0 \quad \forall x \in \mathbb{R}^n \\ A = A^T \end{cases}$$

Theorem [eigenvalues]:

A symmetric matrix A is positive definite iff all its eigenvalues are positive

Partial ordering of square matrices:

For arbitrary square matrices M, N we write $M \geq N$

if $M - N \geq 0$; i.e., $M - N$ is positive semi-definite.

Example: PSD matrices

Theorem [Cone of PSD matrices]:

The set of symmetric, PSD matrices form a convex cone:

$$S_+^n = \{A \in \mathbb{R}^{n \times n} \mid A \succeq 0\}$$

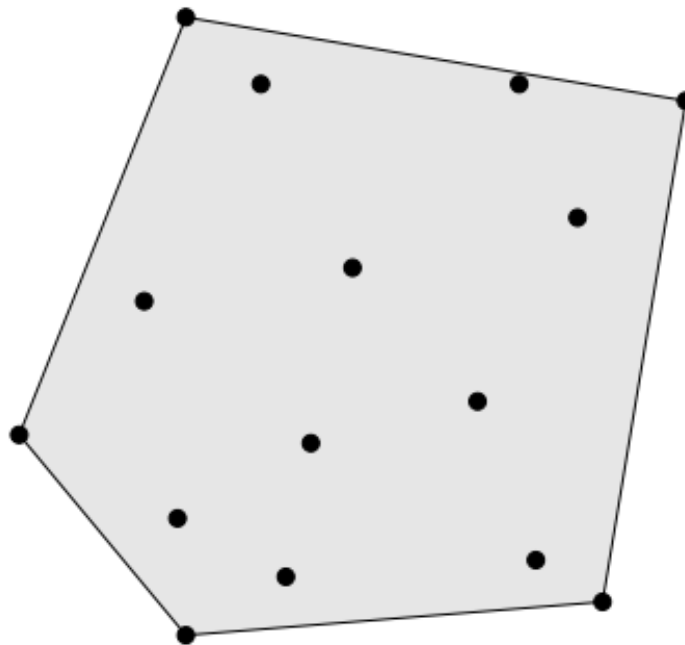
Proof:

$$A, B \in S_+^n \Rightarrow x^T [\theta A + (1-\theta) B] x = \underbrace{\theta \underbrace{x^T A x}_{\geq 0} + (1-\theta) \underbrace{x^T B x}_{\geq 0}}_{\geq 0}$$

Convex set representation with convex hull

Theorem: [Representation of a closed convex set with a convex hull]

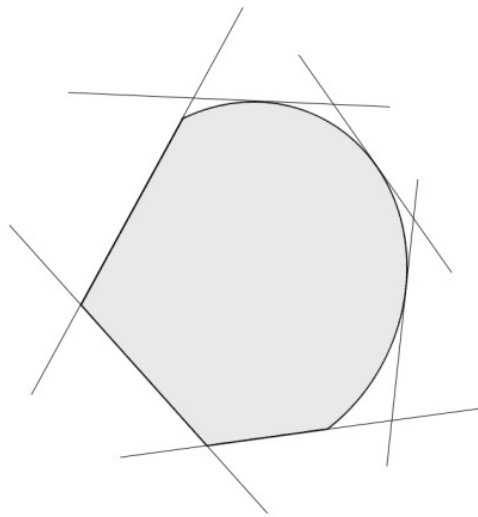
LET $S \subseteq \mathbb{R}^n \Rightarrow C = \mathcal{C}(\text{CONV } S)$ IS THE
CONVEX HULL OF POSSIBLY INFINITE
MANY POINTS IN S



Dual representation

Theorem: [Representation of a closed convex set with half spaces]

LET $S \subseteq \mathbb{R}^n$. $\Rightarrow \overbrace{\mathcal{Q}(\text{CONV } S)}^C$ IS THE INTERSECTION
OF ALL THE CLOSED HALF-SPACES WHICH CONTAIN S
 $\{x \mid a_i^T x + b_i \leq 0\}$



Convexity-preserving set operations

□ Translation $C + b$

□ Scaling αC

□ Intersection $C, D \text{ convex} \Rightarrow C \cap D \text{ convex}$

CAN BE EXTENDED TO INFINITE NUMBER OF SETS
IF S_α IS A CONVEX SET $\forall \alpha \in A \Rightarrow \bigcap_{\alpha \in A} S_\alpha$ IS CONVEX

Convexity-preserving set operations

□ Affine function

- E.g. projection, dropping coordinates

$$\text{If } C \subset \mathbb{R}^n \text{ CONVEX, } A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^m \\ \Rightarrow AC + b = \{Ax + b \mid x \in C\} \subset \mathbb{R}^m \text{ is CONVEX.}$$

□ Set sum

$$C_1 + C_2 = \{c_1 + c_2 \mid c_1 \in C_1, c_2 \in C_2\}$$

□ Direct sum

$$C_1 \times C_2 = \{(c_1, c_2) \in \mathbb{R}^{n+m}, c_1 \in C_1, c_2 \in C_2\}$$

Convex Optimization

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5. Convexity Part II

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MACHINE LEARNING DEPARTMENT



Convex Hull

Definition 1

$$\text{conv}[C] = \left\{ \bigcap_{\alpha} S_{\alpha} \mid S_{\alpha} \subset \mathbb{R}^n, C \subset S_{\alpha}, S_{\alpha} \text{ convex} \right\}$$

Note:

- If C is a finite set, then this is closed polyhedron.
- If C contains infinite many points, then this can be open, closed, or none of them

Theorem [Definition 2, Primal representation]

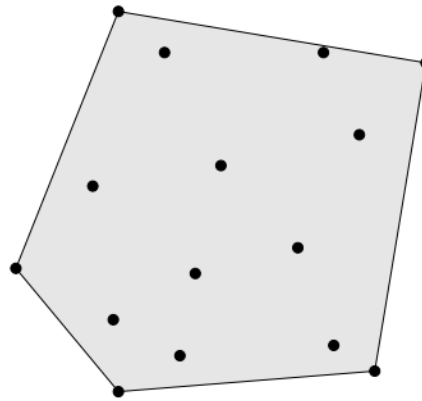
$$\text{conv}[C] = \left\{ \theta_1 x_1 + \dots + \theta_k x_k \mid x_i \in C, \theta_i \geq 0 \forall i=1 \dots k, \sum_{i=1}^k \theta_i = 1, k \in \mathbb{Z}_+ \right\}$$

A closed convex set is the intersection of all the closed half spaces containing S

Convex set representation with convex hull

Theorem: [Representation of a closed convex set with a convex hull]

LET $S \subseteq \mathbb{R}^n \Rightarrow C = \text{CH}(S)$ IS THE
CONVEX HULL OF POSSIBLY INFINITE
MANY POINT IN S



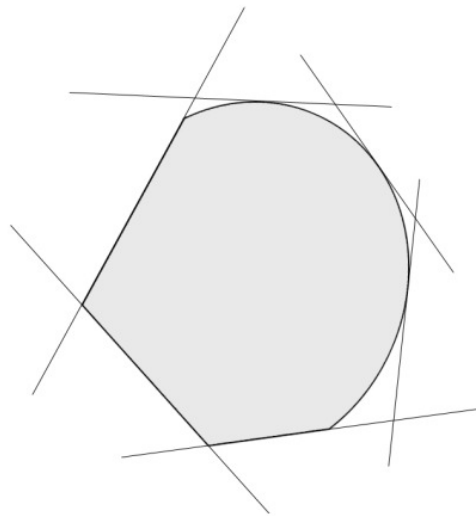
$$\text{CONV}[C] = \left\{ \theta_1 x_1 + \dots + \theta_k x_k \mid x_i \in C, \theta_i \geq 0 \forall i=1 \dots k, \sum_{i=1}^k \theta_i = 1, k \in \mathbb{Z}_+ \right\}$$

Convex hull = convex combination of possibly infinite many points in the set.

Dual representation

Theorem: [Representation of a closed convex set with half spaces]

LET $S \subseteq \mathbb{R}^n$. $\Rightarrow \overbrace{\mathcal{C}(\text{CONV } S)}^C$ IS THE INTERSECTION
OF ALL THE CLOSED HALF-SPACES WHICH CONTAIN S
 $\{x \mid a_i^T x + b_i \leq 0\}$



A closed convex set is the intersection of all the closed half spaces containing S

Convexity-preserving set operations

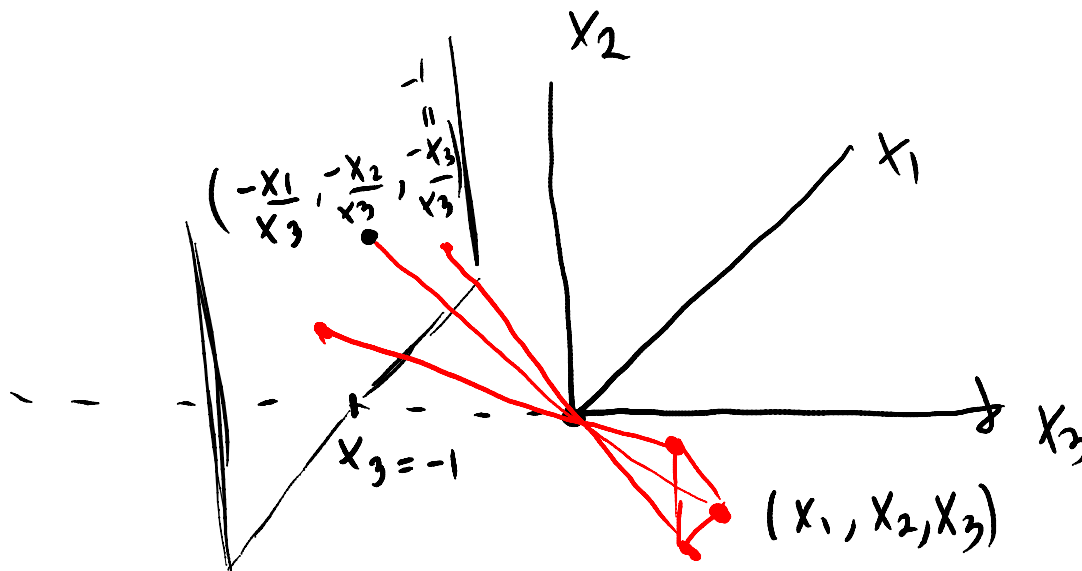
□ Perspective projection (pinhole camera)

IF $C \subset \mathbb{R}^n \times \mathbb{R}_{++}$ IS A CONVEX SET, THEN

$P(C)$ IS ALSO CONVEX

$$P(x) = P(\underbrace{x_1, x_2, \dots, x_n}_{\bar{x}}, t) = (x_1/t, x_2/t, \dots, x_n/t) \in \mathbb{R}^n$$

$$P(\bar{z}, t) = \frac{\bar{z}}{t}$$



Convexity-preserving set operations

□ Linear-fractional function

(perspective function with affine function)

$$f(x) = \frac{Ax + b}{c^T x + d} \quad \begin{array}{ll} A \in \mathbb{R}^{m \times n} & c \in \mathbb{R}^n \\ b \in \mathbb{R}^m & d \in \mathbb{R} \end{array}$$

$$\text{dom } f = \{x \mid c^T x + d > 0\}$$

Theorem: [Image of Linear fractional function]

$$C \subset \mathbb{R}^n \text{ is CONVEX} \Rightarrow f(C) \subset \mathbb{R}^m \text{ is CONVEX}$$

Convexity-preserving set operations

Application: [Conditional probabilities]

$$C = \{ X = (p_{11}, p_{12}, \dots, p_{mn}) \in [0, 1]^{m \times n} \}$$

$U \in \{1, 2, \dots, m\}$ DISCRETE

$V \in \{1, 2, \dots, n\}$ RANDOM
VARIABLES

$$p_{ij} = \text{PROB}(U=i, V=j)$$

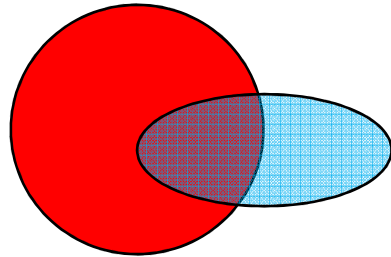
$$f_{ij} = \frac{p_{ij}}{\sum_{k=1}^n p_{kj}} = \text{PROB}(U=i | V=j)$$

$$D = \{ y = (f_{11}, f_{12}, \dots, f_{mn}) \in [0, 1]^{m \times n} \}$$

IF C IS CONVEX $\Rightarrow D$ IS CONVEX

Convexity-preserving set operations

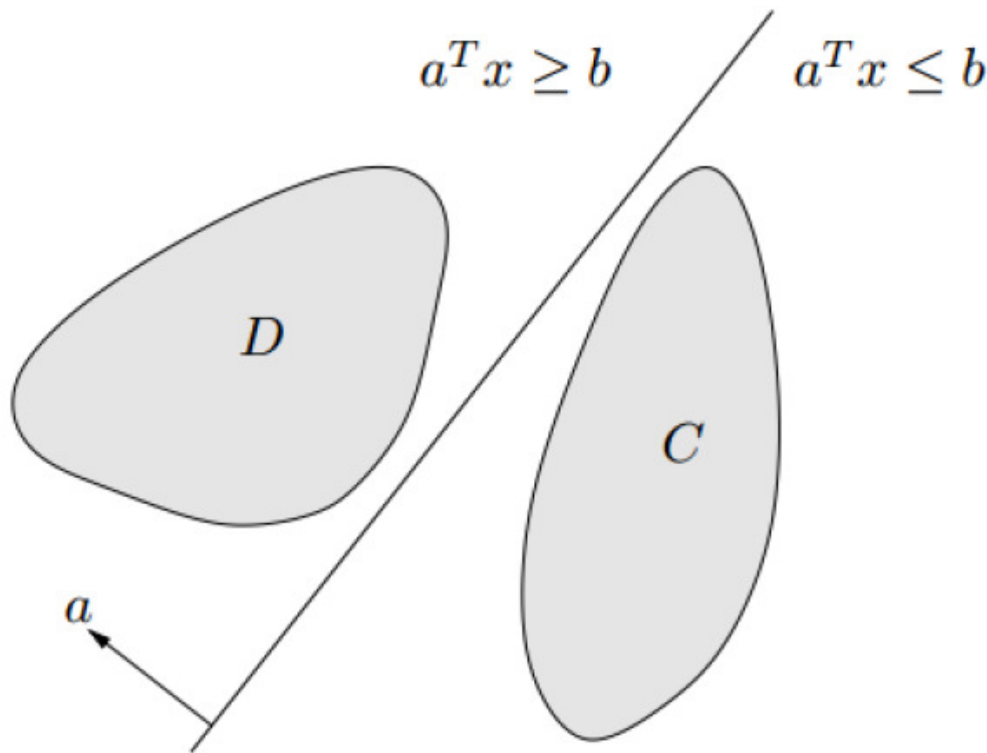
- Union doesn't preserve convexity



Separating hyperplane thm

Theorem: [Separating hyperplane theorem]

$$\left. \begin{array}{l} C, D \text{ CONVEX SETS} \\ C \cap D = \emptyset \end{array} \right\} \Rightarrow \exists \begin{array}{l} a \in \mathbb{R}^n \\ b \in \mathbb{R} \end{array} \text{ s.t. } \begin{array}{l} \forall x \in C \quad a^T x \leq b \\ \forall y \in D \quad a^T y \geq b \end{array}$$



Separating hyperplane thm

Definition: [Strong separation]

$$\begin{aligned} a^T [C_1 + B(0, \varepsilon)] &> b \\ a^T [C_2 + B(0, \varepsilon)] &< b \end{aligned}$$

Definition: [Proper separation],

IT'S NOT THE CASE THAT BOTH

$$\begin{aligned} C_1 &\subseteq \{x : a^T x = b\} \\ C_2 &\subseteq \{x : a^T x = b\} \end{aligned}$$

Definition: [Strict separation]

$$\begin{aligned} a^T x &> b \quad x \in C_1 \\ a^T x &< b \quad x \in C_2 \end{aligned}$$

It "strictly separates" them if neither one touches the hyperplane.

Separating hyperplane thm

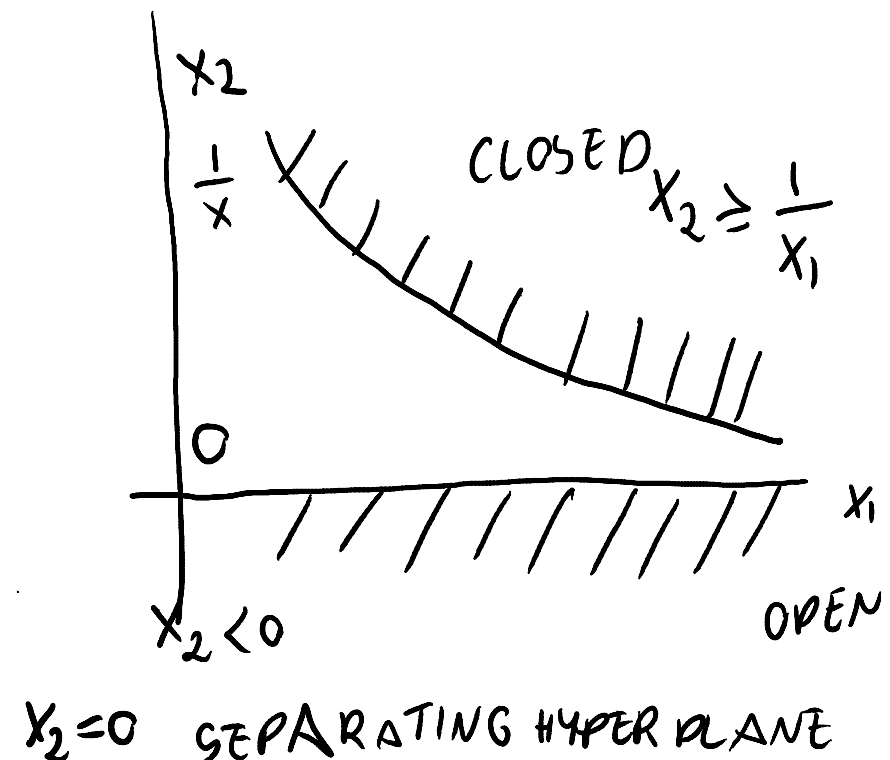
Theorem: [Strong separation theorem]

C_1, C_2 NON EMPTY CONVEX SET IN \mathbb{R}^n
 $\text{cl } C_1 \cap \text{cl } C_2 = \emptyset$
 EITHER C_1 OR C_2 IS BOUNDED

$\Rightarrow \exists$ HYPERPLANE
 SEPARATING C_1 AND C_2
 STRONGLY

Counterexample:

Why do we need at least one bounded set?



Separating hyperplane thm II

Theorem: [Strong separation theorem II]

C_1, C_2 ARE NONEMPTY CONVEX SETS

$C_1 \subseteq \mathbb{R}^n, C_2 \subseteq \mathbb{R}^n$

\exists HYPERPLANE SEPARATING C_1 AND C_2 STRONGLY

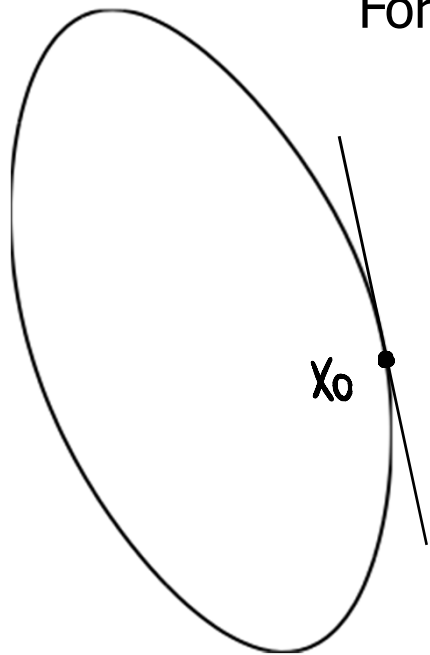
$$\Leftrightarrow \inf_{\substack{x_1 \in C_1 \\ x_2 \in C_2}} \{ |x_1 - x_2| \} > 0$$

$$\Leftrightarrow \text{DIST}(C_1, C_2) > 0$$

$$\Leftrightarrow 0 \in \text{cl}(C_1 - C_2)$$

Supporting hyperplane thm

Theorem: [Supporting hyperplane theorem]



For any point x_0 on the boundary of convex C

$$\exists \text{ HYPERPLANE } \{x \mid a^T x = b\}$$

$$a \neq 0$$

$$\text{s.t. } \forall x \in C \quad a^T x \leq a^T x_0$$

$$a^T x_0 = b$$

Theorem: [Partial converse of the supporting hyperplane theorem]

- SET C IS CLOSED
 - $\text{INT } C \neq \emptyset$
 - $x_0 \in \text{bd } C \Rightarrow \exists \text{ SUPPORTING HYPERPLANE OF } C$
- $\} \Rightarrow C \text{ IS CONVEX}$

Proving a set convex

- ❑ Use definition directly
- ❑ Represent as convex hull
- ❑ Represent as the intersection of halfspaces
- ❑ Supporting hyperplane partial converse:
 - C closed, nonempty interior, has supporting hyperplane at all boundary points $\Rightarrow C$ convex
- ❑ Build C up from simpler sets using convexity-preserving operations

Convex functions

Convex functions

Definition [convex function]:

A function $f: \mathbb{R}^n \rightarrow \mathbb{R}$ is CONVEX IF

- $\text{DOM } f$ IS A CONVEX SET
- $f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$
 $\forall x, y \in \text{DOM } f$
 $\forall \theta \in [0, 1]$

Definition [strictly convex function]:

$f(\theta x + (1-\theta)y) < \theta f(x) + (1-\theta)f(y)$
 $\forall x \neq y$
 $\forall 0 < \theta < 1$

FOR EXAMPLE x^4 IS STRICTLY CONVEX

Concave functions

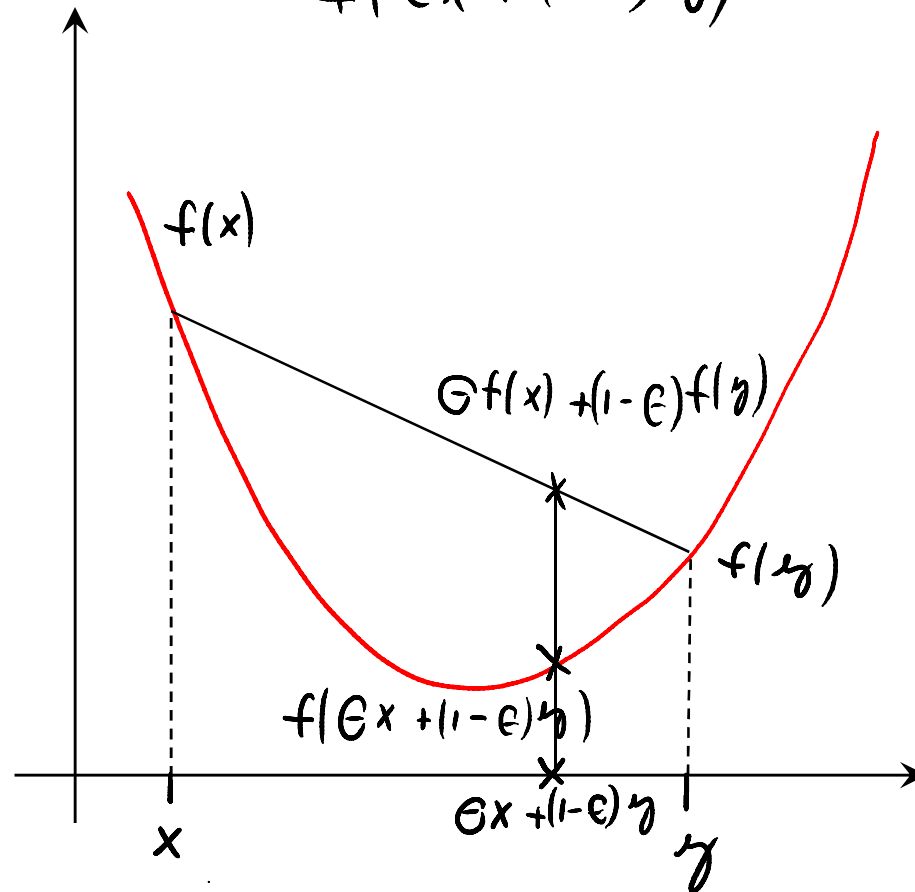
Definition [concave function]:

$-f$ is convex

Convex functions

Geometric interpretation

$$f(\epsilon x + (1-\epsilon)y) \leq \epsilon f(x) + (1-\epsilon)f(y)$$



Strongly convexity

Definition: [m-strongly convex function ($m > 0$)

$$(\nabla f(x) - \nabla f(y))^T (x - y) \geq m \|x - y\|_2^2$$

An equivalent condition:

$$f(y) \geq f(x) + \nabla f(x)^T (y - x) + \frac{m}{2} \|y - x\|_2^2$$

Without gradient: $\forall t \in [0, 1]$

$$f(tx + (1 - t)y) \leq tf(x) + (1 - t)f(y) - \frac{1}{2}mt(1 - t)\|x - y\|_2^2$$

With Hessian:

$$\nabla^2 f(x) \succeq mI \text{ for all } x \text{ in the domain}$$

A strongly convex function is also strictly convex, but not vice-versa.

Examples

$f(x) = x^4$: convex, strictly convex, not strongly convex.

$f(x) = |x|$: convex, not strictly convex.

Examples: Convex functions

- Convex**
- $|x|^p$ $p \geq 1$ ON \mathbb{R}
 - $f(x) = \max(x_1, \dots, x_n)$ ON \mathbb{R}^n
 - ANY NORM

- Concave**
- GEOMETRIC MEAN $f(x) = \left(\prod_{i=1}^n x_i \right)^{1/n}$ CONCAVE ON \mathbb{R}_{++}^n
 - $\log \det(X)$ IS CONCAVE ON S_{++}^n

Extended reals

We can extend f from $\text{dom } f$ to \mathbb{R}^n without changing its convexity

LET $\tilde{f}: \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$

$$\tilde{f}(x) = \begin{cases} f(x) & x \in \text{dom } f \\ \infty & x \notin \text{dom } f \end{cases}$$

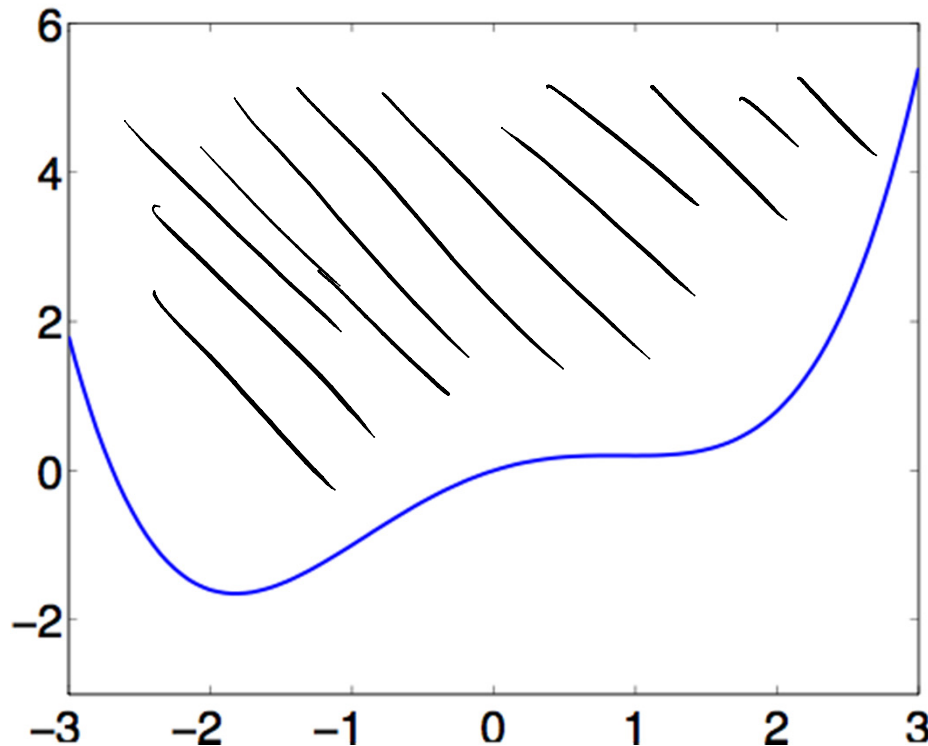
Theorem:

f IS CONVEX $\Leftrightarrow \tilde{f}$ IS CONVEX

$$\Leftrightarrow \tilde{f}(\epsilon x + (1-\epsilon)y) \leq \epsilon \tilde{f}(x) + (1-\epsilon) \tilde{f}(y)$$

Epigraph

Definition [epigraph]: $\text{EPI}(f) = \{ (x, t) : x \in \text{DOM } f, t \geq f(x) \}$



Theorem [convexity of the epigraph]:

$$f: \mathbb{R}^n \rightarrow \mathbb{R} \text{ is convex} \Leftrightarrow \text{EPI}(f) \text{ is convex}$$

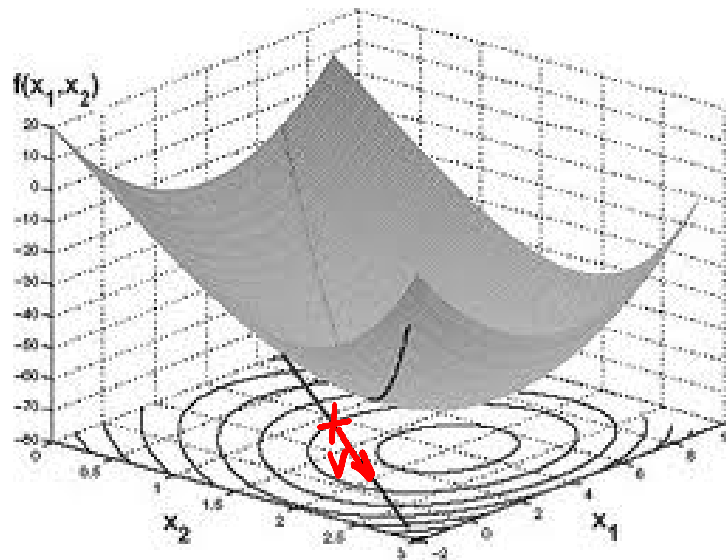
Convex Function Properties

0th order characterization

$$f \text{ CONVEX} \iff g(t) = f(x + tv) \text{ IS CONVEX}$$

$f: \mathbb{R}^n \rightarrow \mathbb{R}$ ON $\text{DOM } g = \{t \mid x + tv \in \text{DOM } f\}$
 $\forall x \in \text{DOM } f \text{ AND } \forall v \in \mathbb{R}^n$

This is useful, because we only need to check the convexity of 1D functions.



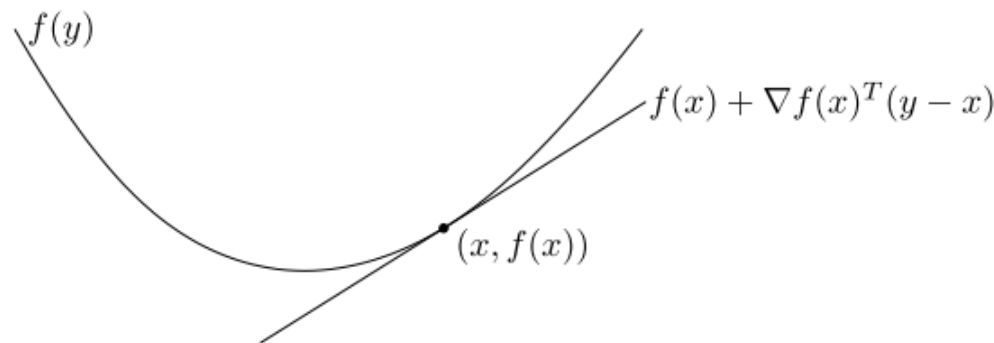
Graph courtesy of Prof. Robert Freund

Convex Function Properties

1st order characterization

- LET f BE A DIFFERENTIABLE FUNCTION
- $\text{DOM } f$ OPEN, CONVEX

$$f \text{ IS CONVEX} \iff f(y) \geq f(x) + \nabla f(x)^T (y - x) \quad \forall y$$



The 1st order Taylor approximation is a global underestimator of f .

$$\text{Corollary: } \left. \begin{array}{l} f \text{ CONVEX} \\ \nabla f(x) = 0 \end{array} \right\} \Rightarrow \begin{array}{l} f(y) \geq f(x) \quad \forall y \\ x \text{ IS GLOBAL OPTIMUM} \end{array}$$

Convex Function Properties

2nd order characterization

- LET f BE TWICE DIFFERENTIABLE
- $\text{DOM } f$ OPEN

$$f \text{ IS CONVEX } \Leftrightarrow \nabla^2 f(x) \succeq 0 \quad \forall x \in \text{DOM } f$$

Lemma IF $\nabla^2 f(x) \succ 0 \quad \forall x \in \text{DOM } f \Rightarrow f \text{ IS STRICTLY CONVEX}$
 \nLeftarrow FOR EXAMPLE x^4

Jensen's inequality

Theorem

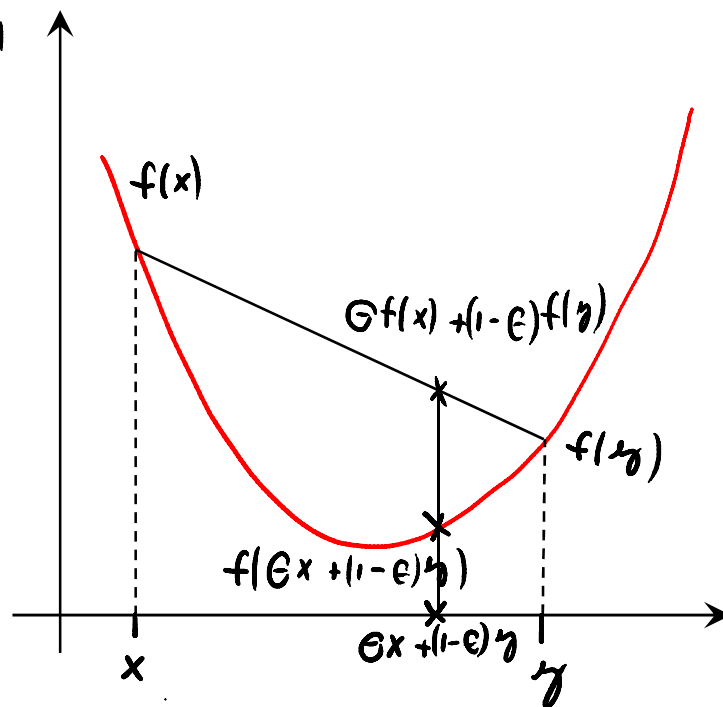
$$f: \mathbb{R}^r \rightarrow (-\infty, \infty]$$

$$f \text{ is convex} \Leftrightarrow f(\lambda_1 x_1 + \dots + \lambda_m x_m) \leq \lambda_1 f(x_1) + \dots + \lambda_m f(x_m)$$

$$\forall \lambda_1 \geq 0, \dots, \lambda_m \geq 0, \sum_{i=1}^m \lambda_i = 1$$

$$f \text{ is convex} \Rightarrow f(\mathbb{E} X) \leq \mathbb{E} f(X)$$

$$X \sim P$$



$$f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$$

Proving a function convex

- ❑ Use definition directly
- ❑ Prove that epigraph is convex via set methods
- ❑ 0^{th} , 1^{st} , 2^{nd} order convexity properties
- ❑ Construct f from simpler convex fns using convexity-preserving ops

Convexity-preserving fn ops

Nonnegative weighted sum

If f_1, f_2 CVX, $w_i \geq 0 \Rightarrow h(x) = w_1 f_1(x) + w_2 f_2(x)$ CVX

Pointwise max/sup

If f, g CVX, $\Rightarrow m(x) = \max\{f(x), g(x)\}$ CVX

Extension of pointwise max/sup

If $f(x, y)$ is convex in x for each y
 $\Rightarrow g(x) = \sup_{y \in C} f(x, y)$ is convex in x ,
provided $g(x) > -\infty$ for some x .

Convexity-preserving fn ops

Affine map

If $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is cvx, $\Rightarrow g(x) = f(Ax + b)$ is cvx,
where $A \in \mathbb{R}^{n \times m}$, $b \in \mathbb{R}^n$.

Composition

If f, g are cvx, and g is non-decreasing,
 $\Rightarrow h(x) = g(f(x))$ is cvx.

If f is concave and g is cvx and non-increasing,
 $\Rightarrow h(x) = g(f(x))$ is cvx.

Perspective map

If $f(x)$ is convex, $\Rightarrow g(x, t) = tf(x/t)$ is convex.

Summary

□ **Convex sets**

- Representation:
 - convex hull, intersect hyperplanes
- supporting, separating hyperplanes
- operations that preserve convexity

□ **Convex functions**

- epigraph
- 0 orders, 1st order, 2nd order conditions
- operations that preserve convexity