

## Lecture 13: KKT conditions

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## 13.1 Continued from last lecture on duality

### 13.1.1 Weak Duality

Recall from last lecture, we reached the following conclusion:

$$f^* = \min_{x \in C} f(x) \geq \min_{x \in C} L(x, u, v) \geq \min_{x \in \mathbb{R}} L(x, u, v) := g(u, v) \quad (13.1)$$

And we denote the tight upper bound of  $g(u, v)$  as  $g^* := \max_{u > 0} g(u, v)$ .

The key insight is that the **weak duality property** always hold no matter the primal problem is convex or not, namely:

$$f^* \geq g^* \quad (13.2)$$

Also note that the dual problem is always a **convex optimization** problem (maximizing a concave function), even when the primal problem is non-convex.

By definition:

$$g(u, v) = \min_{x \in \mathbb{R}^n} \left[ f(x) + \sum_{i=1}^m u_i h_i(x) + \sum_{j=1}^r v_j \ell_j(x) \right] \quad (13.3)$$

$$= - \max_{x \in \mathbb{R}^n} \left[ -f(x) - \sum_{i=1}^m u_i h_i(x) - \sum_{j=1}^r v_j \ell_j(x) \right] \quad (13.4)$$

For any  $x$ , pointwise maximum is a convex function in  $(u, v)$ .

The following example illustrates this property:

$$\begin{aligned} \min_x f(x) &= x^4 - 50x^2 + 100x \\ \text{subject to } x &\geq -4.5 \end{aligned} \quad (13.5)$$

The original problem is obvious non-convex as shown in Fig. 13.1.

Though the dual function can be derived explicitly (differentiate the Lagrangian and find a closed-form solution of a cubic equation), the form of  $g$  is quite complicated but it is concave!

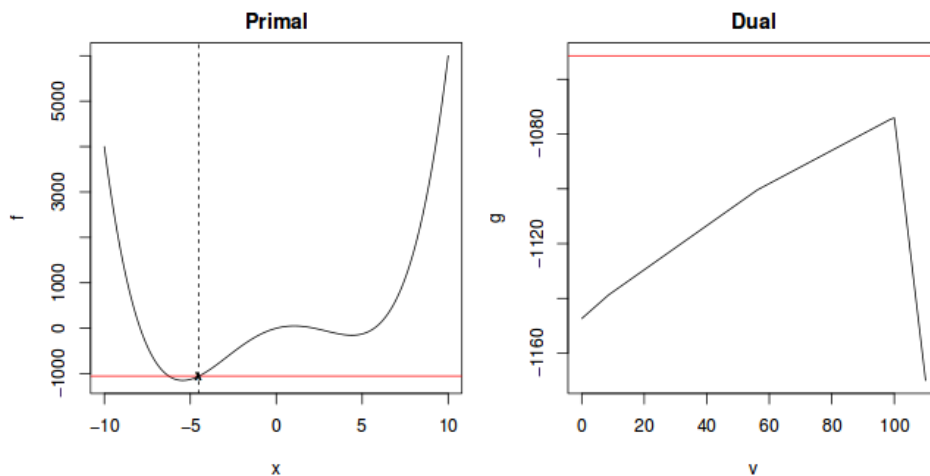


Figure 13.1: Nonconvex primal problem and its concave dual problem

### 13.1.2 Strong Duality

Weak duality is good but in many problems we have observed something even better:

$$f^* = g^* \quad (13.6)$$

which is called the **strong duality**. But when do we have this nice property?

**Slater's Condition:**

- if the primal is convex (i.e.,  $f$  and  $h_1, \dots, h_m$  are convex,  $\ell_1, \dots, \ell_r$  are affine)
- if there exists at least one strictly feasible  $x \in \mathbb{R}^n$   
(i.e.,  $h_1(x) < 0, \dots, h_m(x) < 0$  and  $\ell_1(x) = 0, \dots, \ell_r(x) = 0$ )

This is actually a weak statement and it can be further refined: need strict inequality only over  $h_i$  that are not affine.

In the case of linear programming:

- If the primal LP is feasible, then by Slater's condition strong duality holds and hence  $f^* = g^*$ ;
- If the dual LP is feasible, then by Slater's condition strong duality holds and hence  $g^* = f^*$ ;
- Strong duality breaks only when both primal and dual are infeasible.

## 13.2 Recap and Summary: Primal problem and dual problem

**Primal problem:**

$$\begin{aligned} & \min_{x \in \mathbb{R}^n} f(x) \\ & \text{subject to } h_i(x) \leq 0, i = 1, \dots, m \\ & \quad \quad \ell_j(x) = 0, j = 1, \dots, r \end{aligned} \quad (13.7)$$

**Lagrangian:**

$$L(x, u, v) = f(x) + \sum_{i=1}^m u_i h_i(x) + \sum_{j=1}^r v_j \ell_j(x) \quad (13.8)$$

**Lagrange dual function:**

$$g(u, v) = \min_{x \in \mathbb{R}^n} L(x, u, v) \quad (13.9)$$

**Dual problem:**

$$\begin{aligned} & \max_{u \in \mathbb{R}^m, v \in \mathbb{R}^r} g(u, v) \\ & \text{subject to } u \geq 0. \end{aligned} \quad (13.10)$$

**Immediate results:**

- For any feasible solution,  $f(x) \geq L(x, u, v)$ .
- $g(u, v)$  is always concave, even if  $f(x)$  is not convex.
- Weak duality: It is always true that  $f^* \geq g^*$ . Hence for any  $(x, u, v)$ ,

$$f(x) - f^* \leq f(x) - g(u, v) \quad (13.11)$$

If  $f(x) - g(u, v) = 0$ , then  $x$  is primal optimal and  $(u, v)$  are dual optimal.

- Slater's condition: for convex primal, if there is an  $x$  such that

$$h_1(x) < 0, \dots, h_m(x) < 0 \text{ and } \ell_1(x) = 0, \dots, \ell_r(x) = 0 \quad (13.12)$$

then strong duality holds, i.e.,  $f^* = g^*$ .

### 13.3 Karush-Kuhn-Tucker conditions

**Theorem 13.1** Under strong duality,  $x^*$  and  $u^*, v^*$  are primal and dual solutions if and only if the KKT conditions hold, which are:

- **Stationarity:**  $0 \in \partial f(x^*) + \sum_{i=1}^m u_i^* \partial h_i(x^*) + \sum_{j=1}^r v_j^* \partial \ell_j(x^*)$
- **Complementary slackness:**  $u_i^* h_i(x^*) = 0$  for all  $i$
- **Primal feasibility:**  $h_i(x^*) \leq 0, \ell_j(x^*) = 0$  for all  $i, j$
- **Dual feasibility:**  $u_i^* \geq 0$  for all  $i$ .

**Proof:** We first prove necessity.

$$f(x^*) = g(u^*, v^*) \quad (13.13)$$

$$= \min_{x \in \mathbb{R}^n} f(x) + \sum_{i=1}^m u_i^* h_i(x) + \sum_{j=1}^r v_j^* \ell_j(x) \quad (13.14)$$

$$\leq \min_{x \in \mathbb{R}^n} f(x) \quad (13.15)$$

$$= f(x^*) \quad (13.16)$$

Hence the above inequality is actually an equality, which means

- Primal feasibility and dual feasibility obviously hold;
- $x^*$  minimizes  $L(x, u^*, v^*)$  over  $\mathbb{R}^n$ , hence the subdifferential of  $L(x, u^*, v^*)$  contains 0 at  $x = x^*$ , which is the stationarity condition;
- $\sum_{i=1}^m u_i^* h_i(x^*) = 0$  and since  $u_i^* \leq 0$  and  $h_i(x) \leq 0$ , hence  $u_i^* h_i(x^*) = 0$  for all  $i$ , which is the complementary slackness.

Next we prove sufficiency. If there exists  $x^*$ ,  $u^*$  and  $v^*$  that satisfy the KKT condition, then

$$g(u^*, v^*) = \min_{x \in \mathbb{R}^n} L(x, u^*, v^*) \quad (13.17)$$

$$= f(x^*) + \sum_{i=1}^m u_i^* h_i(x^*) + \sum_{j=1}^r v_j^* \ell_j(x^*) \quad (\text{stationarity}) \quad (13.18)$$

$$= f(x^*) \quad (\text{complementary slackness, dual feasibility}) \quad (13.19)$$

which means the duality gap is zero and therefore  $x^*$ ,  $u^*$ ,  $v^*$  are optimal solutions. ■

**Warning:** One may attempt to conclude that stationarity is equivalent to the following:

$$0 = \nabla f(x) + \sum_{i=1}^m u_i^* \nabla h_i(x) + \sum_{j=1}^r v_j^* \nabla \ell_j(x) \quad (13.20)$$

This is only true when  $f(x)$ ,  $h_i(x)$  and  $\ell_j(x)$  are convex.

Another way to formulate the problem is using the indicator function  $\mathbb{I}$  and normal cone  $\mathcal{N}$ :

$$f(x) + \sum_{i=1}^m \mathbb{I}_{\{h_i(x) \leq 0\}} + \sum_{j=1}^r \mathbb{I}_{\{\ell_j(x) = 0\}} \quad (13.21)$$

$$0 \in \partial f(x^*) + \sum_{i=1}^m \mathcal{N}_{\{h_i(x^*) \leq 0\}} + \sum_{j=1}^r \mathcal{N}_{\{\ell_j(x^*) = 0\}} \quad (13.22)$$

$$\iff x^* \text{ is optimal} \quad (13.23)$$

## 13.4 Examples

### 13.4.1 Quadratic optimization with equality constraints

Consider for  $Q \succeq 0$ ,

$$\begin{aligned} \min_{x \in \mathbb{R}^n} \quad & \frac{1}{2} x^T Q x + c^T x \\ \text{subject to} \quad & Ax = 0 \end{aligned} \tag{13.24}$$

As  $Q \succeq 0$ , the above problem is convex. By stationarity and primal feasibility, we have  $x$  is a solution if and only if

$$\begin{bmatrix} Q & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} x \\ v \end{bmatrix} = \begin{bmatrix} -c \\ 0 \end{bmatrix} \tag{13.25}$$

for some  $v$ .

### 13.4.2 Side note: Newton's method on linearly constrained problem

$$\min f(x) \tag{13.26}$$

$$\text{subject to } Ax = b \tag{13.27}$$

Recall Newton's method updates  $x$  as follows:  $x^+ = x - (\nabla^2 f(x))^{-1} \nabla f(x)$ , but this update rule cannot guarantee the constraint is satisfied at each step.

Instead we can start from one  $x$  that satisfies the constraint:  $Ax = b$ , then we update it with:

$$x^+ = x + \Delta \quad \text{where } A \Delta = 0 \tag{13.28}$$

and then minimize the following quadratic problem:

$$\frac{1}{2} \Delta^T (\nabla^2 f(x)) \Delta + \nabla f(x)^T \Delta \tag{13.29}$$

### 13.4.3 Water-filling

Consider

$$\begin{aligned} \min_{x \in \mathbb{R}^n} \quad & - \sum_{i=1}^n \log(\alpha_i + x_i) \\ \text{subject to} \quad & x \geq 0, 1^T x = 1 \end{aligned} \tag{13.30}$$

The Lagrangian is

$$L(x, u, v) = - \sum_{i=1}^n \log(\alpha_i + x_i) - \sum_{i=1}^n u_i x_i + v \left( \sum_{i=1}^n x_i - 1 \right)$$

Stationarity:

$$-\frac{1}{\alpha_i + x_i} - u_i + v = 0 \quad \text{for all } i = 1, \dots, n \quad (13.31)$$

Complementary slackness:

$$u_i x_i = 0 \quad \text{for all } i = 1, \dots, n \quad (13.32)$$

Primal feasibility:

$$x \geq 0, \mathbf{1}^T x = 1 \quad (13.33)$$

Dual feasibility:

$$u_i \geq 0 \quad (13.34)$$

Combining the above results, we get

$$\begin{aligned} v - \frac{1}{\alpha_i + x_i} &\geq 0 \\ x_i(v - \frac{1}{\alpha_i + x_i}) &= 0 \end{aligned} \quad (13.35)$$

Hence if  $v < \frac{1}{\alpha_i}$ , then  $x_i > 0$ , then  $v = \frac{1}{\alpha_i + x_i}$  which is  $x_i = \frac{1}{v} - \alpha_i$ ;

if  $v \geq \frac{1}{\alpha_i}$ , then  $x_i = 0$ . In sum, we get  $x_i = \max\{0, \frac{1}{v} - \alpha_i\}$ . By primal feasibility, we solve the uni-variate optimization problem

$$\sum_{i=1}^n \max\{0, \frac{1}{v} - \alpha_i\} = 1 \quad (13.36)$$

to get the solution to the original problem.

#### 13.4.4 Lasso

Consider

$$\min_{\beta \in \mathbb{R}^p} \frac{1}{2} \|y - X\beta\|_2^2 + \lambda \|\beta\|_1 \quad (13.37)$$

From stationarity,

$$X^T(y - X\beta) = \lambda s \quad (13.38)$$

where  $s \in \partial \|\beta\|_1$ , that is

$$s_i \in \begin{cases} \text{sign}(\beta_i) & \text{if } \beta_i \neq 0 \\ [-1, 1] & \text{if } \beta_i = 0 \end{cases} \quad (13.39)$$

from which we directly get if  $|X_i^T(y - X\beta)| < \lambda$ , then  $\beta_i = 0$ .

### 13.4.5 Group Lasso

Consider

$$\min_{\beta=(\beta^{(1)},\beta^{(2)},\dots,\beta^{(G)})\in\mathbb{R}^p} \frac{1}{2}\|y - X\beta\|_2^2 + \lambda \sum_{i=1}^G w_i \|\beta^{(i)}\|_2 \quad (13.40)$$

From stationarity, for  $i = 1, 2, \dots, G$ ,

$$(X^{(i)})^T(y - X\beta) = \lambda w_i s^{(i)} \quad (13.41)$$

$$\text{where } s^{(i)} \in \partial\|\beta^{(i)}\|_2 = \begin{cases} \frac{\beta^{(i)}}{\|\beta^{(i)}\|_2} & \text{if } \beta^{(i)} \neq 0 \\ \{v : \|v\|_2 \leq 1\} & \text{otherwise} \end{cases}.$$

- If  $\|(X^{(i)})^T(y - X\beta)\|_2 < \lambda w_i \Rightarrow \beta^{(i)} = 0$ .
- If  $\beta^{(i)} \neq 0$ ,

$$\begin{aligned} (X^{(i)})^T(y - X^{(i)}\beta^{(i)} - \sum_{j \neq i} X^{(j)}\beta^{(j)}) &= \lambda w_i \frac{\beta^{(i)}}{\|\beta^{(i)}\|_2} \\ \Rightarrow -(X^{(i)})^T X^{(i)}\beta^{(i)} + (X^{(i)})^T r^{(i)} &= \lambda w_i \frac{\beta^{(i)}}{\|\beta^{(i)}\|_2} \\ \Rightarrow \left( \lambda w_i \frac{\beta^{(i)}}{\|\beta^{(i)}\|_2} I + (X^{(i)})^T X^{(i)} \right) \beta^{(i)} &= (X^{(i)})^T r^{(i)} \\ \Rightarrow \beta^{(i)} &= \left( \lambda w_i \frac{\beta^{(i)}}{\|\beta^{(i)}\|_2} I + (X^{(i)})^T X^{(i)} \right)^{-1} (X^{(i)})^T r^{(i)} \end{aligned} \quad (13.42)$$

$$\text{where } r^{(i)} = y - \sum_{j \neq i} X^{(j)}\beta^{(j)}.$$

## 13.5 Summary

Under strong duality, we can characterize the primal solution from its dual problem.

Recall that under strong duality, the KKT conditions are necessary for optimality. Given dual solutions  $(u^*, v^*)$ , any primal solution satisfies the stationarity condition:

$$0 \in \partial f(x^*) + \sum_{i=1}^m u_i^* \partial h_i(x^*) + \sum_{j=1}^r v_j^* \partial \ell_j(x^*) \quad (13.43)$$

In other words,  $x^*$  achieves the minimum in  $\min_{x \in \mathbb{R}^n} L(x, u^*, v^*)$ .

- In general, this reveals a characterization of primal solutions
- In particular, if this is satisfied uniquely (i.e., above problem has a unique minimizer), then the corresponding point must be the primal solution.

## References

[BV04] S. BOYD and L. VANDENBERGHE, “Convex optimization”.