

Primal

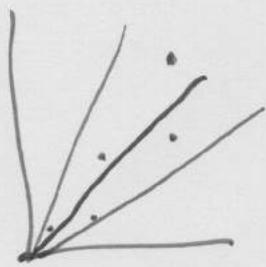
$$\{x : Ax = b, x \geq 0\}$$

$$c^\top$$



$$\begin{aligned} & \uparrow b \\ & \{y, s : \\ & A^\top y + s = c \\ & s \geq 0\} \end{aligned}$$

$$\begin{aligned} \mathcal{F}^0 &\cong \mathbb{R}_{++}^n && \text{bijection} \\ (x, y, s) &\mapsto x_{S1} \end{aligned}$$



Recall central path eqns

$$\begin{bmatrix} A^\top y + s = c \\ Ax = b \\ x_{S1} = \tau \mathbf{1} \\ x, s \geq 0 \end{bmatrix}$$

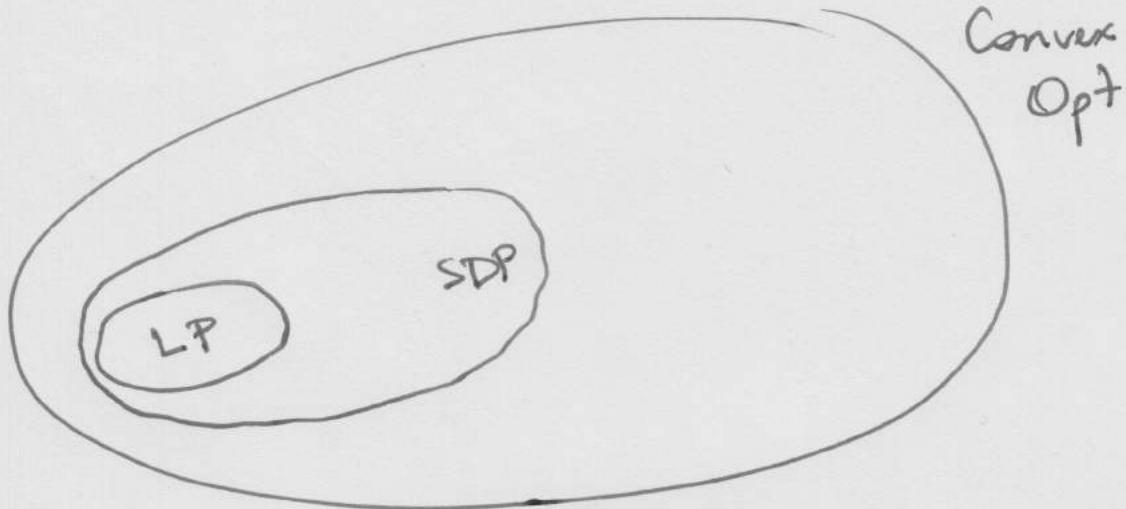
$$\mathbb{R}_{++}^n$$

Newton's method:

optimization:  $\min f(x)$ ,  $x \mapsto x + \Delta x$   
 where  $[\nabla^2 f(x)] \Delta x = -\nabla f(x)$

solving eqns:  $G(x) = 0$ ,  $x \mapsto x + \Delta x$

$$G'(x) \Delta x = -G(x)$$



Primal SDP

$$\begin{aligned} \min \quad & C \cdot X \\ \text{such that} \quad & A(X) = b \\ & X \succeq 0 \end{aligned}$$

$$A : S^n \rightarrow \mathbb{R}^m$$

$$A^* : \mathbb{R}^m \rightarrow S^n$$

$$\text{Typically } A(X) = \begin{bmatrix} A_1 \cdot X \\ \vdots \\ A_m \cdot X \end{bmatrix}$$

$$\rightsquigarrow A^*(y) = \sum_{i=1}^m y_i A_i$$

$$\min 2x_{12} \quad \Leftrightarrow \quad \begin{bmatrix} 0 & x_{12} \\ x_{12} & x_{22} \end{bmatrix} \succcurlyeq 0$$

$$\min c \cdot x \quad A_1 \cdot x = b$$

$$x \succcurlyeq 0$$

for  $c = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$ ,  $A_1 = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$ ,  $b = 0$

Dual  $\max 0 \cdot y$

$$\begin{bmatrix} y & 0 \\ 0 & 0 \end{bmatrix} \succcurlyeq \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 & x_{12} \\ x_{12} & x_{22} \end{bmatrix} \succcurlyeq 0 \Rightarrow x_{12} = 0$$

so primal opt value is 0.

Dual constraint:

$$\begin{bmatrix} -y & 1 \\ 1 & 0 \end{bmatrix} \succcurlyeq 0$$

$y$  can never happen  
no matter what  $y$  is.

$$\min x_{11} \quad \min C \cdot x$$

$$\begin{bmatrix} x_{11} & 1 \\ 1 & x_{22} \end{bmatrix} x \geq 0 \Leftrightarrow \begin{array}{l} A_1 \cdot x = b \\ x \geq 0 \end{array}$$

$$\text{for } C = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, A_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, b = 2$$

$$\text{Dual} \quad \max 2y$$

$$\begin{bmatrix} 0 & y \\ y & 0 \end{bmatrix} \leq \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$$

dual opt value is 0 attained at  $y=0$

Primal problem: opt value is 0 but it is not attained

$$\min a x_{22}$$

$\leftarrow$  opt val = 1  
attained

$$\begin{bmatrix} 0 & x_{12} & 1-x_{22} \\ x_{12} & x_{22} & x_{23} \\ 1-x_{22} & x_{23} & x_{33} \end{bmatrix} x \geq 0$$

$$\text{Dual: } \max 2y_2$$

$\leftarrow$  opt val = 0  
attained

$$\begin{bmatrix} y_1 & 0 & y_2 \\ 0 & 2y_2 & 0 \\ y_2 & 0 & 0 \end{bmatrix} \leq \begin{bmatrix} 0 & 0 & 0 \\ 0 & a & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

Barrier method for primal SDP:

$$\min C \cdot X$$

$$A(X) = b$$

$$X \succcurlyeq 0$$

$$\min C \cdot X - \tau \log \det X$$

$$A(X) = b$$

The function  $F: S_{++}^n \rightarrow \mathbb{R}$

$$X \mapsto -\log \det X$$

"log. barrier function for  $S_{++}^n$ "

Observe

$$\begin{aligned} F(X) &= -\log \det X = -\log \prod_{j=1}^n \lambda_j(X) \\ &= -\sum_{j=1}^n \log(\lambda_j(X)) \end{aligned}$$

cf log barrier for  $\mathbb{R}_+^n$ :  $-\sum \log x_i$

Recall  $f(t) = -\log t \rightsquigarrow f'(t) = -\frac{1}{t}$

For  $F(X) = -\log \det X \rightsquigarrow \nabla F(X) = -X^{-1}$

$d_F(X, S)$  = measure of proximity of  
 $X, S$  to  $\mu I$

A "feasible" method would satisfy

$$A(y) + S - C = 0$$

$$R(x) - b = 0$$

$$XS + \tau I$$

$$G(x, y, S) - \begin{pmatrix} 0 \\ 0 \\ \tau I \end{pmatrix} = 0.$$

"natural" linearization of  $XS - \tau I = 0$

$$\rightsquigarrow X \Delta S + S \Delta X = \tau I - XS$$

Neat calculus fact:  $F(X) = -\log \det X$   
 $\nabla F(X) = -X^{-1}$   
 $\nabla^2 F(X)[Z] = X' Z X'$

Recall Lovasz theta function

$$G = (V, E)$$

$$\Theta(G) = \max_{\substack{\Pi \in \mathbb{R}^{V \times V} \\ \Pi \cdot \Pi^T = I \\ \Pi_{ij} = 0 \quad i, j \in E \\ \Pi \geq 0}} \Pi \cdot \Pi^T \cdot X$$